

Adventures In Real-Time Performance Tuning, Part 1

The real-time for Linux patchset does not guarantee adequate real-time behavior for all target platforms. When using real-time Linux on a new platform you should expect to have to tune the kernel and drivers to provide performance that matches your specific requirements.

This paper provides an example of the trials and tribulations of the tuning journey for a MIPS target board. A brief back of the envelope real-time performance characterization of the board will also be presented.

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Adventures In Real-Time Performance Tuning, Part 1

What I'm trying to tune

Some examples of using the available tuning tools

If I talk fast enough, some performance data for a
MIPS TX4937 processor

What is Real Time?

It is determinism (being able to respond to a stimulus before a deadline) within a given system load envelope.

It is NOT fast response time.

The specific real time application deadlines determine how short the maximum response time must be to deliver real time behavior.

Some examples of deadlines are one second, one millisecond, or five microseconds.

RT latency is the delay from stimulus to when the RT “application” is executing code

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Possible RT application contexts include

- driver interrupt context
- driver thread context
- kernel thread context
- user space thread context

Some components that may contribute to RT latency

- IRQ disabled time
- preempt disabled time
- IRQ latency, from event until bottom half
- RT driver bottom half(s) execution
- non-RT driver bottom half(s) execution
- task switch time

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- IRQ disabled time
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- IRQ latency, from event until bottom half
- RT driver bottom half(s) execution
- non-RT driver bottom half(s) execution
- task switch time

The components that add up to RT latency are important to the tuning process, but keep in mind the end goal of tuning actual RT latency.

Examples of Tuning Knobs

Which hardware is enabled and used

Which kernel functionality is enabled and used

Which drivers are used

Kernel config options

Softirq handling in thread context

Driver top half algorithms

Driver top half polling (vs irq)

Driver bottom half in thread context

Driver bottom half thread priorities

Real-time thread priorities

Non-RT application load

Examples of Tuning Knobs

Kernel algorithms and code

Locks

Timers

CPU affinity and partitioning

- drivers
- kernel threads
- user processes and threads

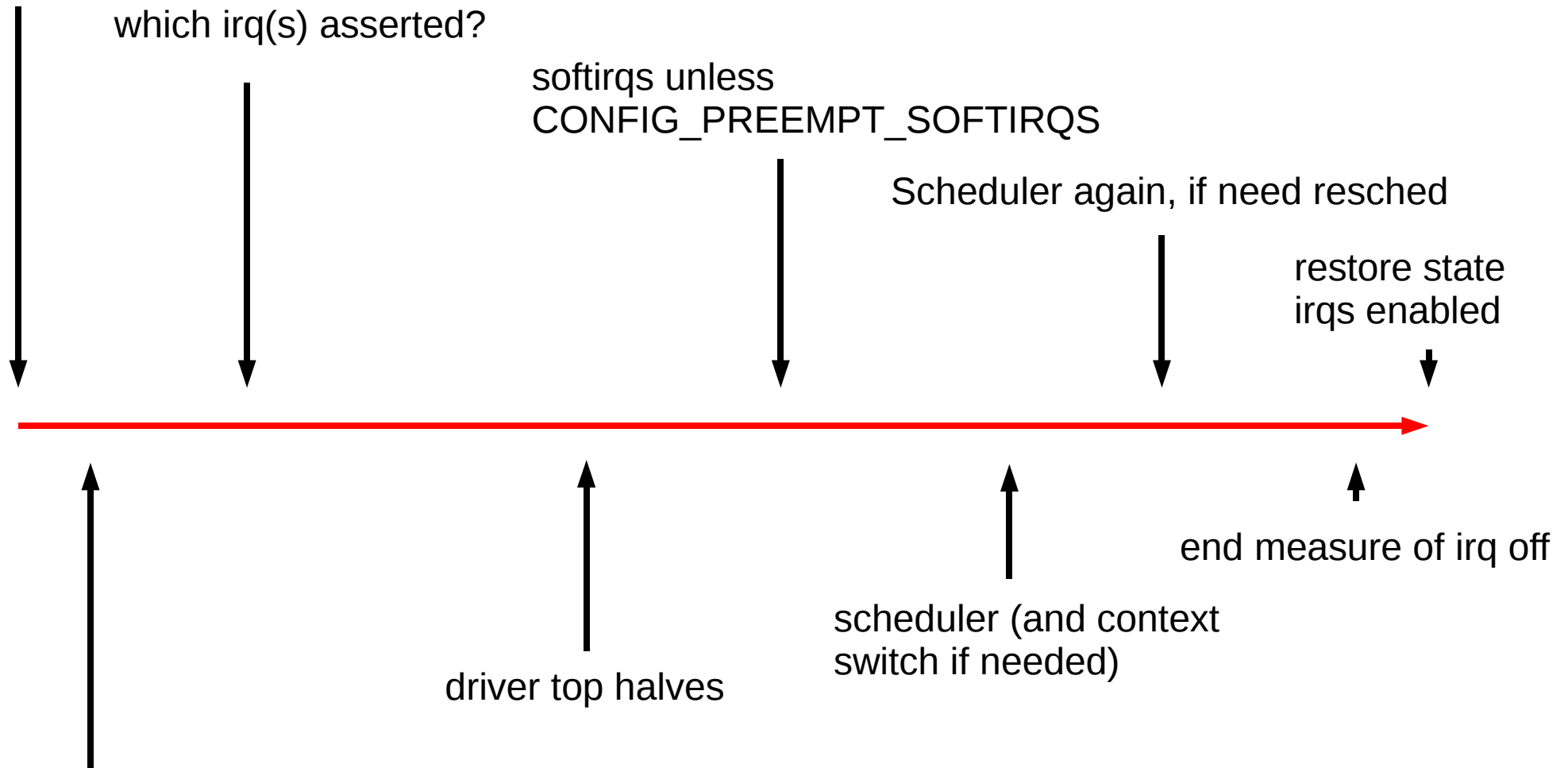
Lock code and data in memory

Lock tlb entries

Lock code and data in cache or fast
memory

Sample Timeline: IRQ Off Latency

interrupt is asserted
irqs disabled



Each arrow points to the completion of the work described by the label.

A Roadmap of my Journey

- 1) add some RT pieces for MIPS and the tx4937 processor
- 2) add MIPS support to RT instrumentation
- 3) tuning
- 4) implement “lite” irq disabled instrumentation
- 5) tuning

Caveats

Tools, instrumentation, techniques, etc are very dependent upon the version of the kernel and rt-preempt patchset.

Kernel data structures, algorithms, performance hot spots change.

An example of instrumentation change:

RT tracing mechanism was accepted in mainline in 2.6.27-rc1. The review process prompted a partial re-write.

To follow this, see LKML:

From: Ingo Molnar

Date: Fri Feb 08 2008

Subject: [git pull] latency tracer

“Linus, please pull the latency tracer tree”

Date: Fri Feb 10 2008

Subject: [00/19] latency tracer

The new tracer, “ftrace”, is also in the rt patchset, starting with patch 2.6.24-rt2:

- cleaned up code
- /debugfs/tracing/ instead of /proc with better control interface
- simultaneous trace of irq off and preempt off
- simultaneous histogram and trace

The examples in this presentation are from the old tracer (mips 2.6.24 + rt patch 2.6.24-rt1), but the ftrace reports look much like the old tracer.

Tuning, part 1

Instrumentation anomalies led to large reported IRQ off times:

- “interrupts disabled” in idle wait (processor specific)
- timer comparison code not handling clock rollover
- timer comparison and capture code not handling switch between raw and non-raw clock sources

The first hint of large latency

```
/proc/latency_hist/interrupt_off_latency/CPU0
```

```
#Minimum latency: 2 microseconds.
```

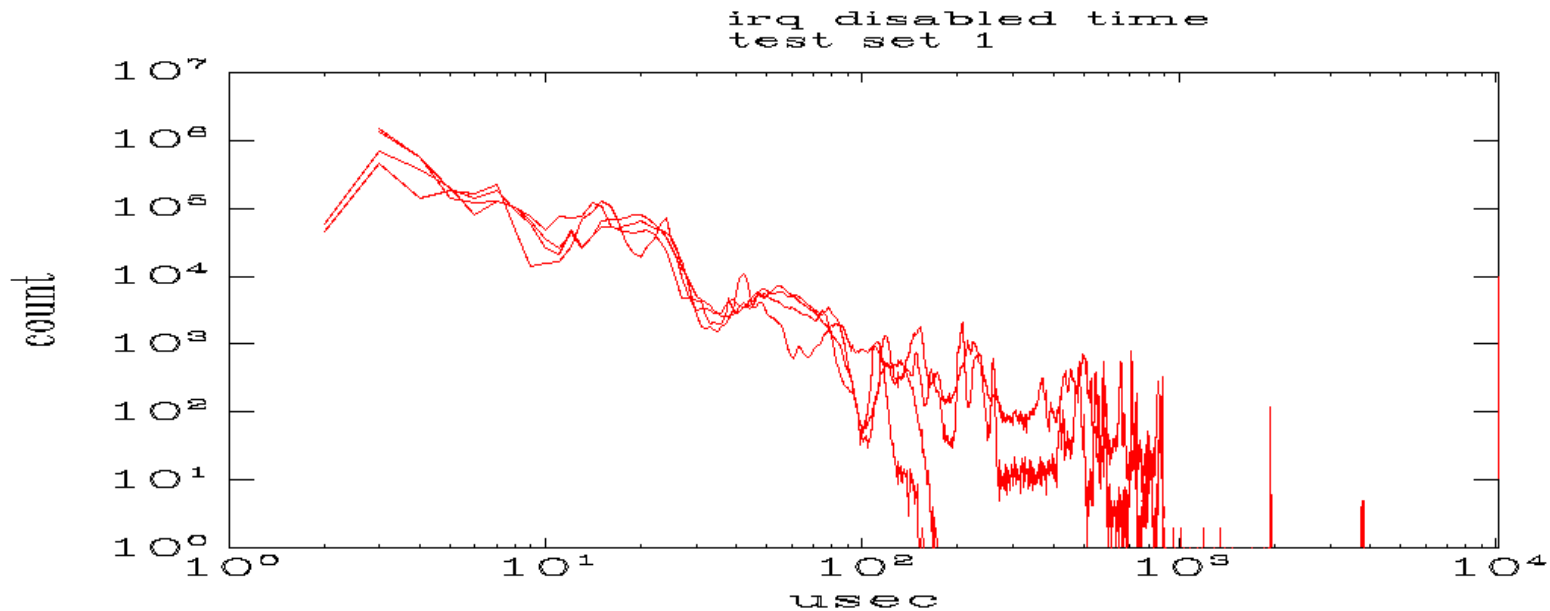
```
#Average latency: 27 microseconds.
```

```
#Maximum latency: 5725129 microseconds.
```

```
#Total samples: 2846758
```

```
#There are 3 samples greater or equal than 10240  
microseconds
```

#usecs	samples
0	0
1	0
2	59063
3	666520
4	362079

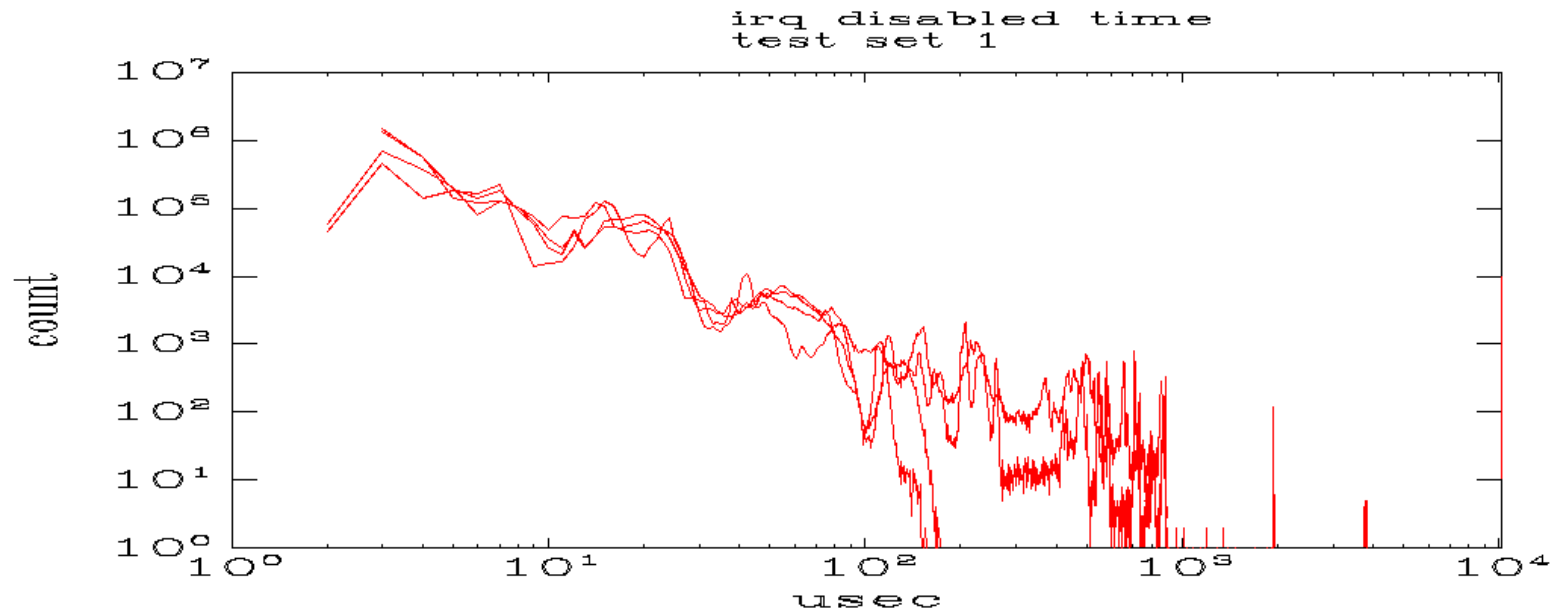


Graphing the histograms provides a quick overview of the data.

Multiple test runs are shown on this graph.

Each line is the result of a single test run.

The vertical bar on the right edge of the graph warns of one or more measurements beyond the right edge of the graph.



There are many irq disabled periods in the range of 200 – 1000 microseconds.

There are a significant number of events beyond 1000 microseconds.

/proc/latency_trace points to a cause

latency: 137 us, #10/10, CPU#0 | (M:rt VP:0, KP:0, S

| task: swapper-0 (uid:0 nice:0 policy:0 rt_prio

=> started at: r4k_wait_irqoff+0x40/0x98 <800213d8>

=> ended at: irq_exit+0xc0/0xf4 <800567c0>

		_-----=> CPU#				
		/ _-----=> irq-s-off				
		/ _-----=> need-resched				
		/ _-----=> hardirq/softirq				
		/ _--=> preempt-depth				
		/				
		delay				
cmd \	/ pid		time \		caller /	

“interrupts disabled” in idle wait

```
switch (c->cputype)
case CPU_TX49XX:
    cpu_wait = r4k_wait_irqoff
```

```
static void r4k_wait_irqoff(void)
{
    local_irq_disable();
    if (!need_resched())
        __asm__("wait \n");
    local_irq_enable();
}
```

The WAIT instruction

The WAIT instruction causes the processor to enter a low-power mode. The processor exits from the low-power mode upon an interrupt exception.

So when an interrupt occurs, `r4k_wait_irqoff()` will immediately re-enable interrupts.

Quick “FIX”

Use the pre-existing MIPS “nowait” boot option

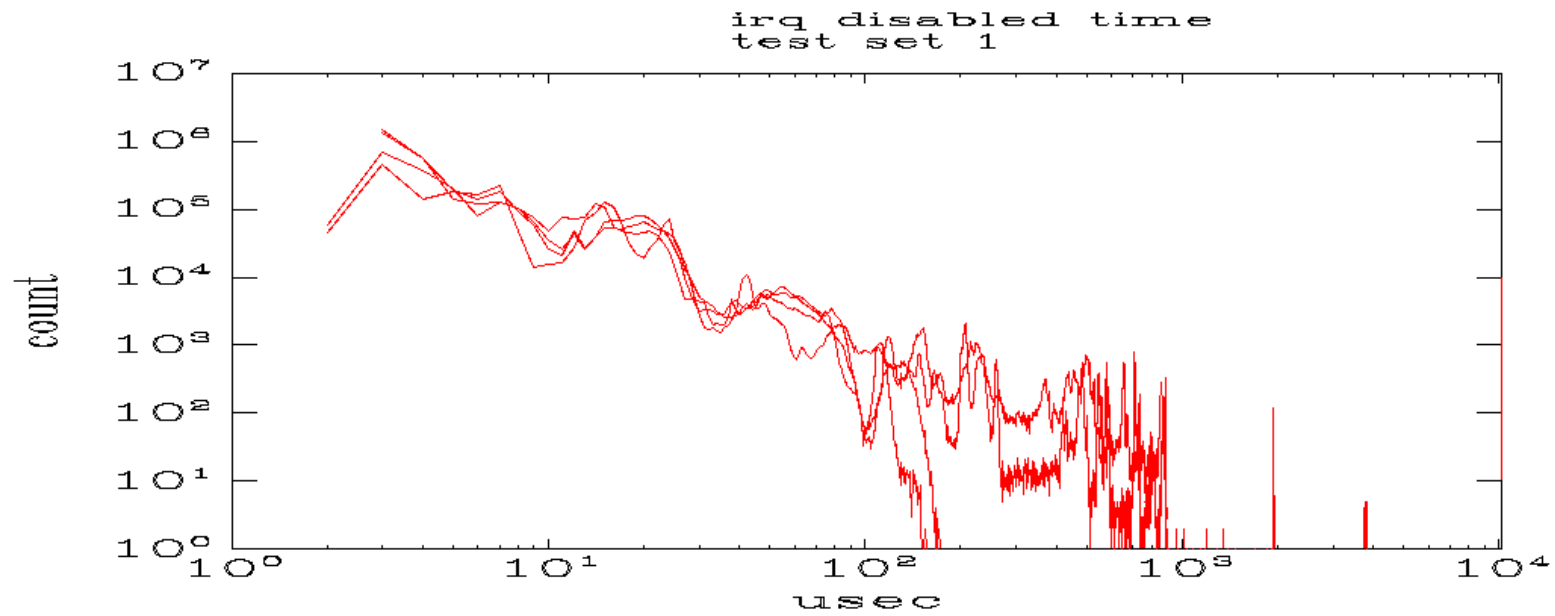
```
static inline void check_wait(void)
{
    if (nowait) {
        printk("Wait instruction disabled.\n");
        return;
    }
    switch (c->cputype)
    case CPU_TX49XX:
        cpu_wait = r4k_wait_irqoff
```

Real “FIX”

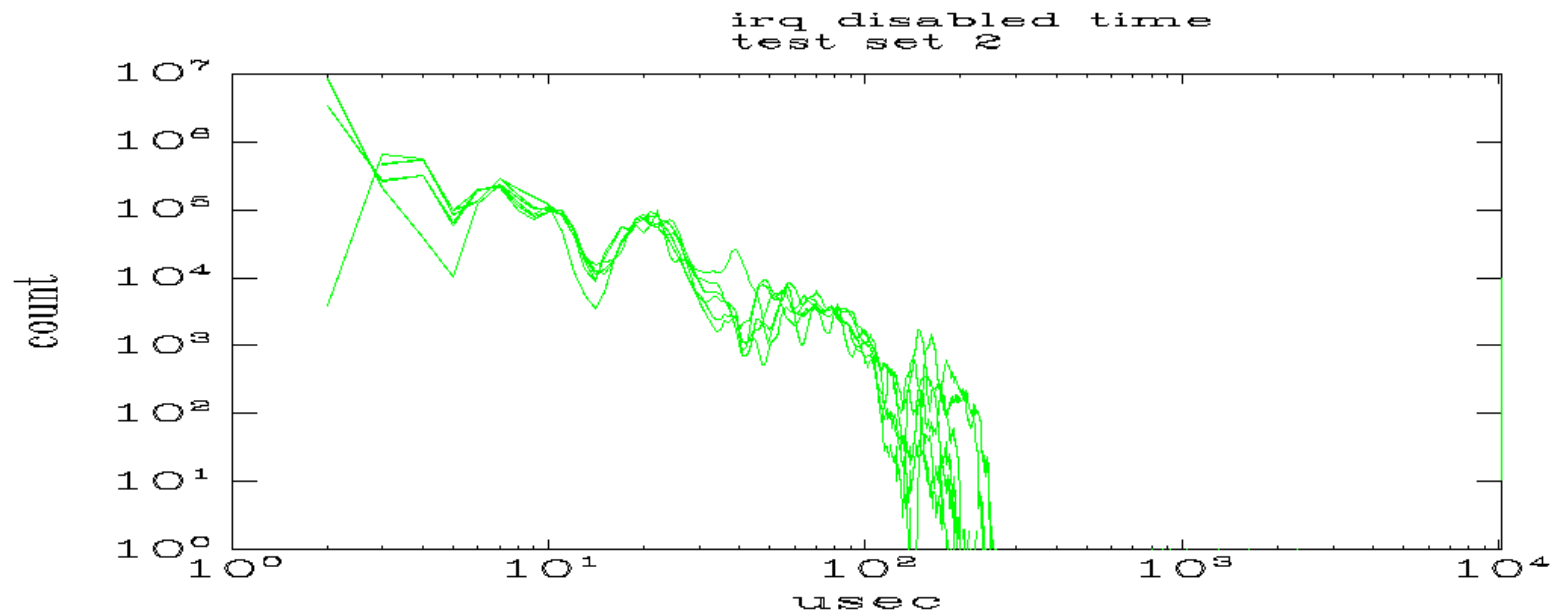
Explicitly stop latency tracing in `cpu_idle()`

For an example, see
`arch/x86/kernel/process_32.c`

Look for
`stop_critical_timing()`
`touch_critical_timing()`
`trace_preempt_enter_idle()`
`trace_preempt_exit_idle()`



Fix cpu_wait()



Very large max latency remains

/proc/latency_hist/interrupt_off_latency/CPU0

#Maximum latency: 5725051 microseconds.

The cause of this was found via normal debugging.

I used the MIPS cycle counter to implement `get_cycles()`, which the latency tracer uses when the raw cycles mode is enabled. The cycle counter holds a 32 bit value, which rolls over quickly. The latency tracer was not coded to handle timer rollover.

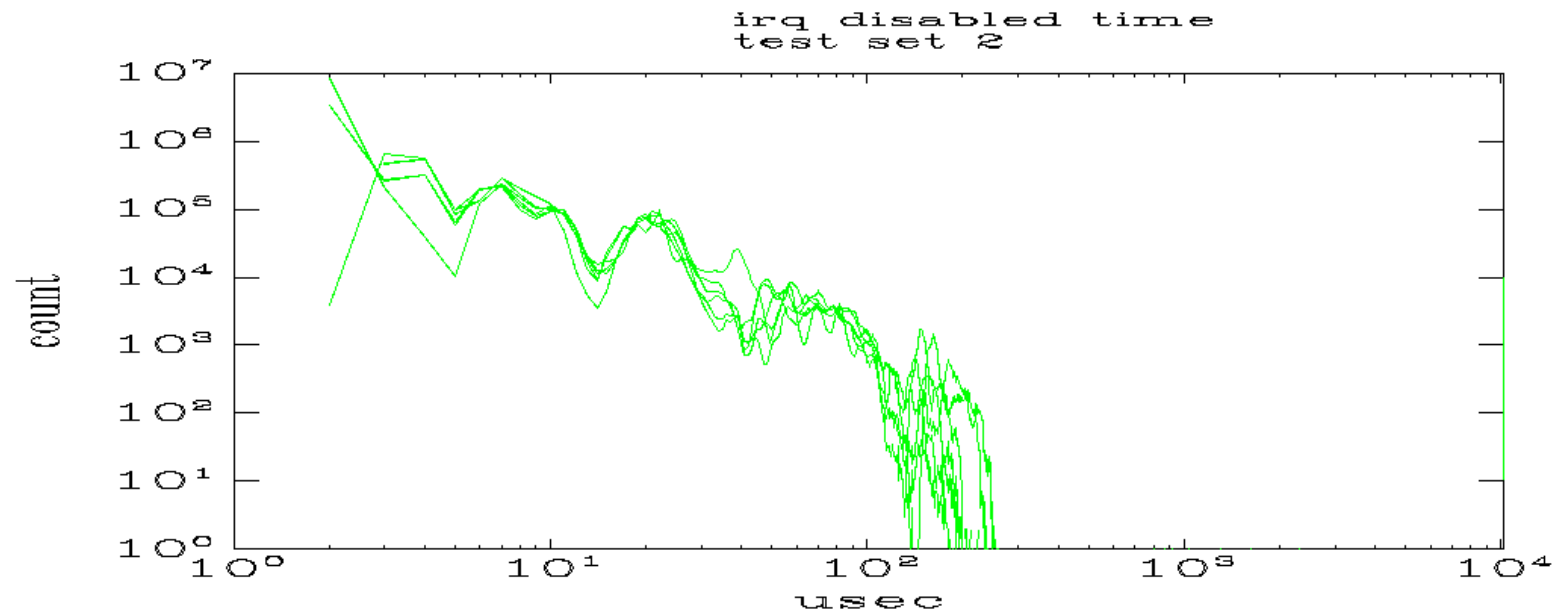
clock rollover FIX

use the same algorithms used for jiffies in
`include/linux/jiffies.h`

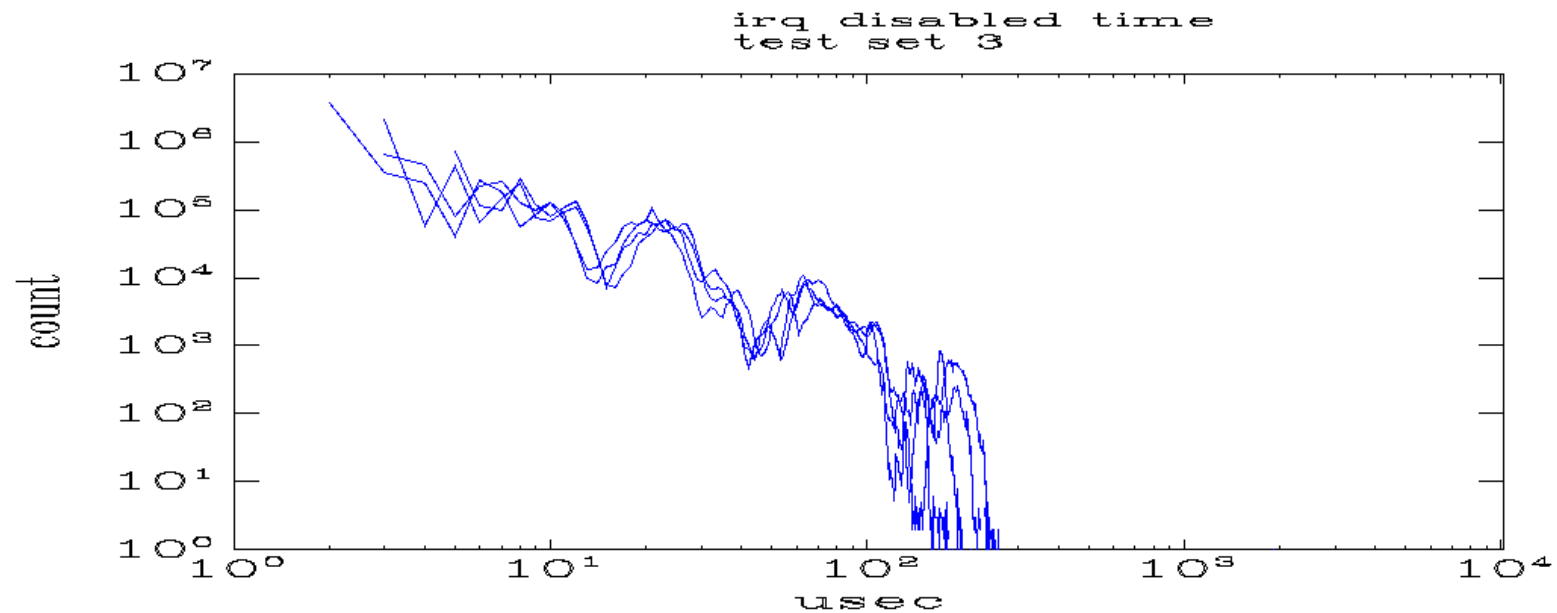
see:

`time_after()`

`time_before()`



Fix timestamp compare



Still large max latency remains

/proc/latency_trace reports large latency

/proc/latency_hist/interrupt_off_latency/CPU0

#Maximum latency: 6765 microseconds.

The cause of this was once again found via normal debugging, not through the RT instrumentation.

switch between raw and non-raw clock sources WORKAROUND

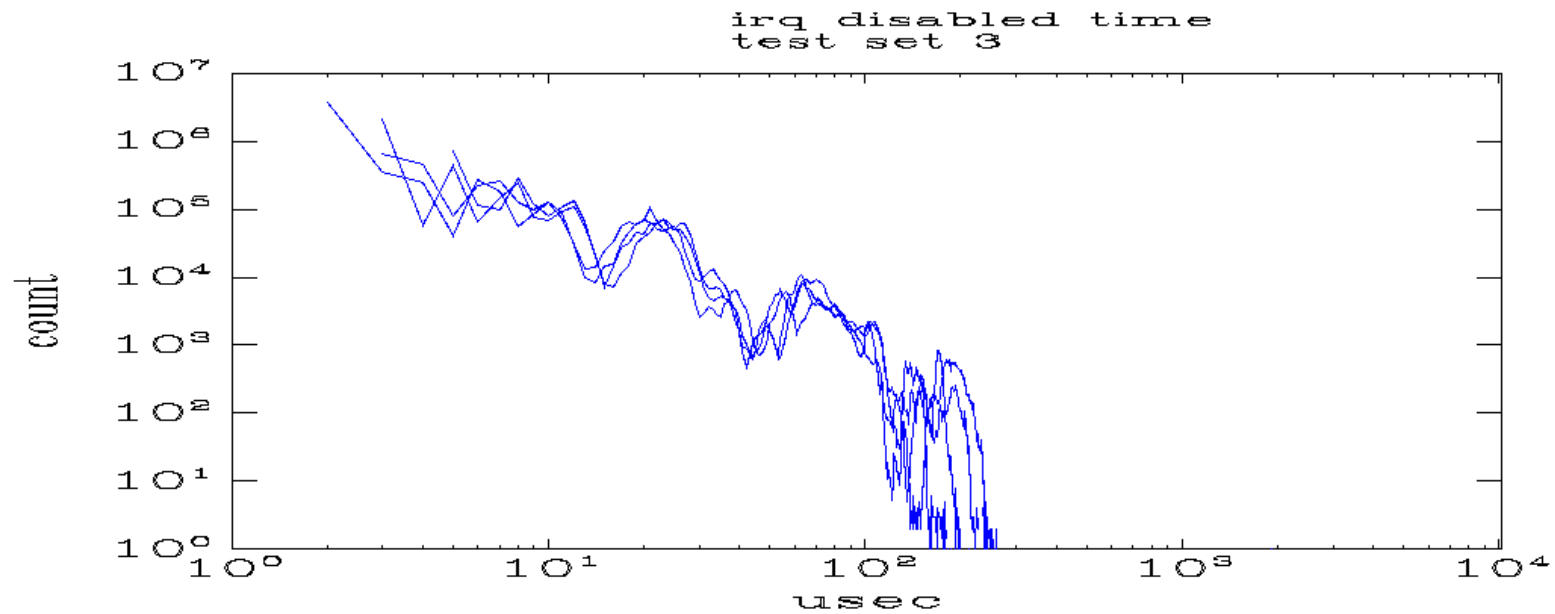
Do not switch back and forth. Use either raw or non-raw for all tracing.

switch between raw and non-raw clock sources FIX

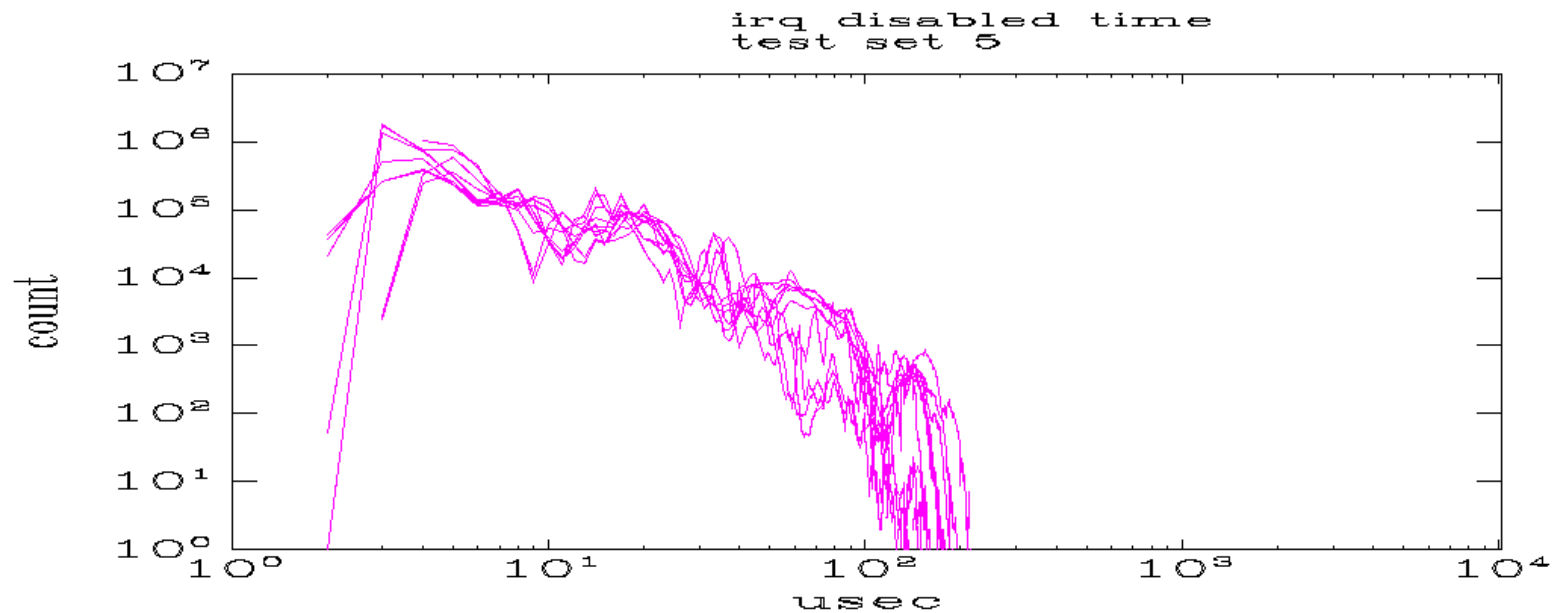
In kernel/latency_trace.c: _____trace()

check for switch between raw and non-raw

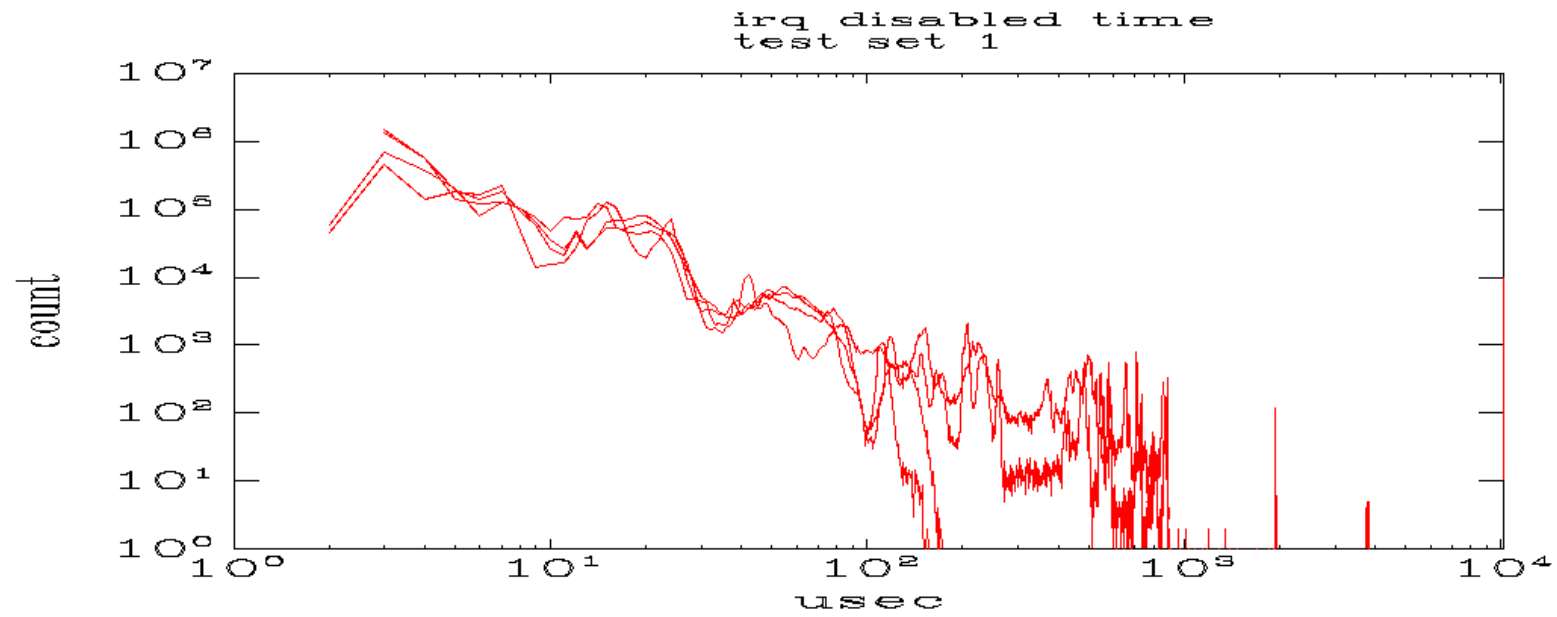
delete timestamps in other mode from current event



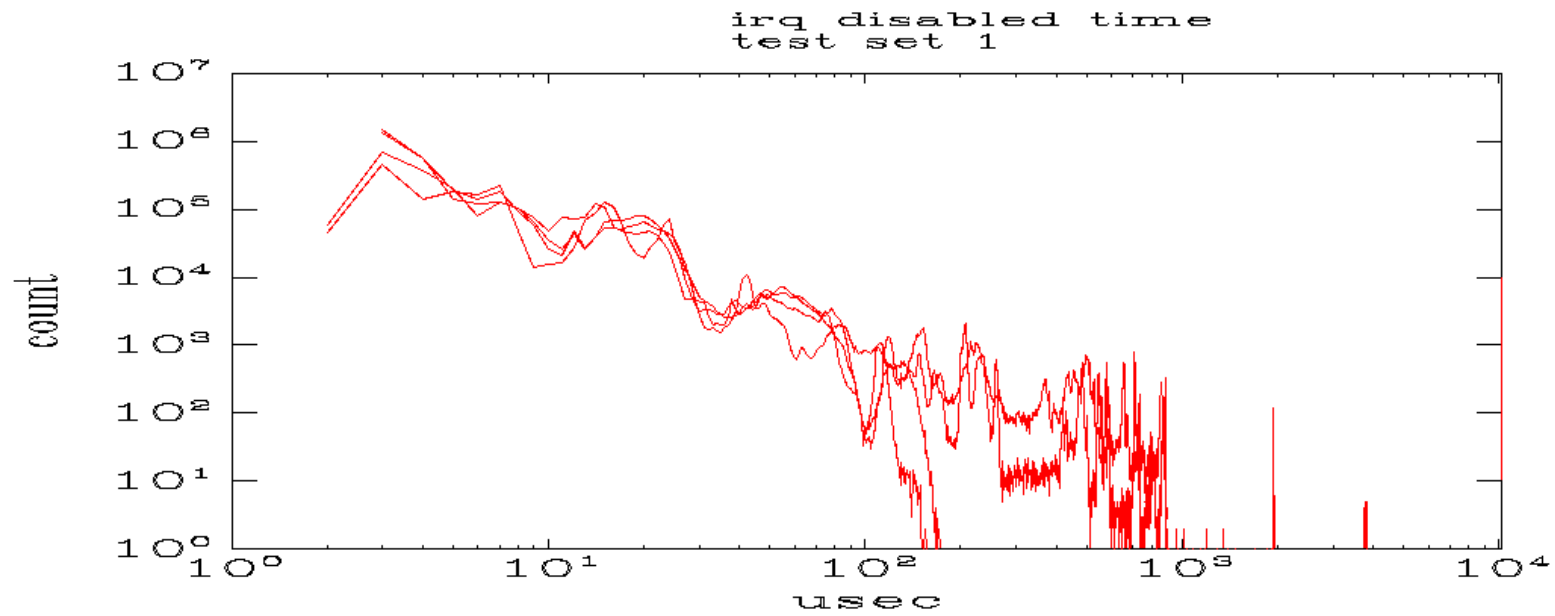
Fix raw / non-raw timestamp transition



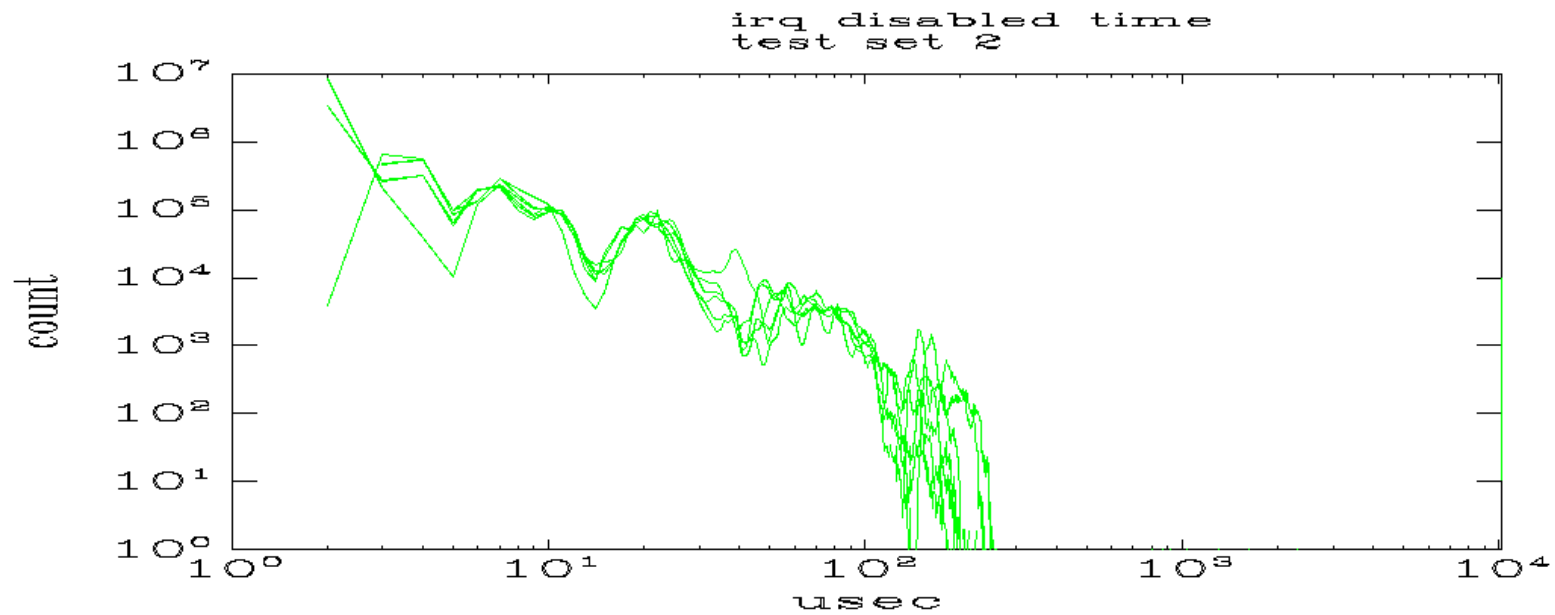
Review of the graphical results

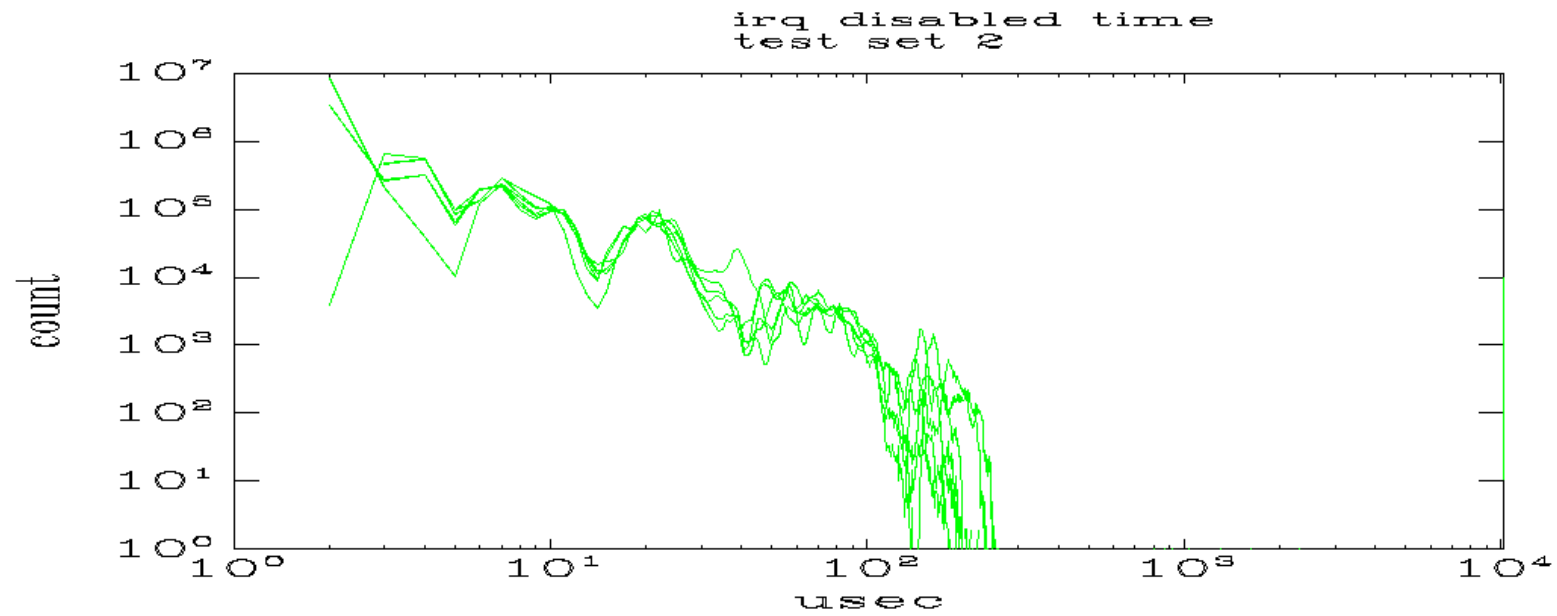


Base measurement

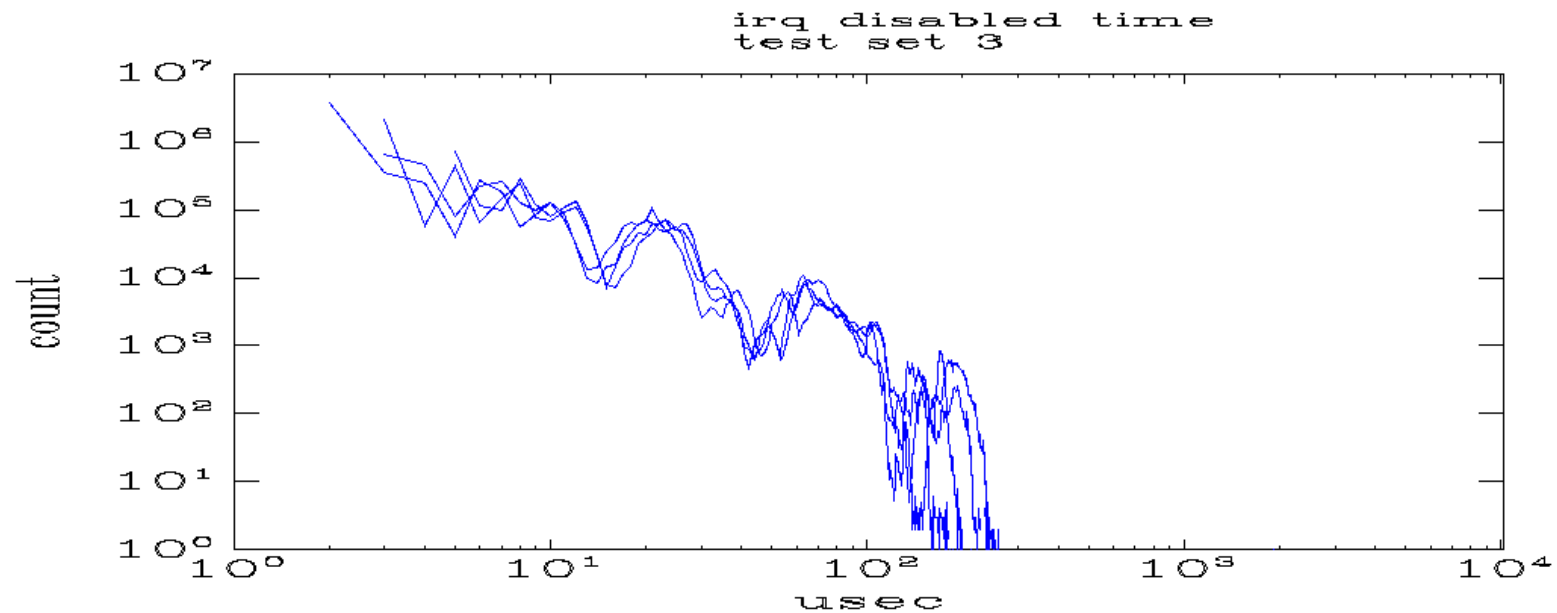


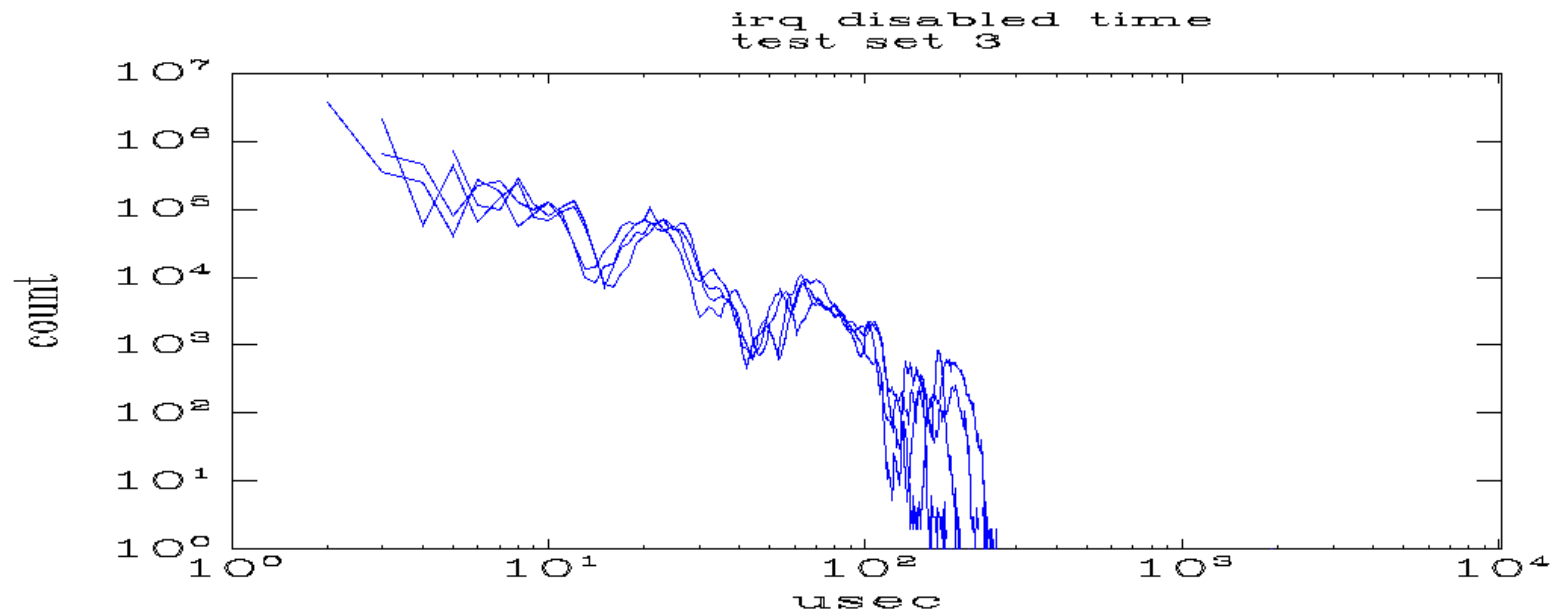
Fix cpu_wait()



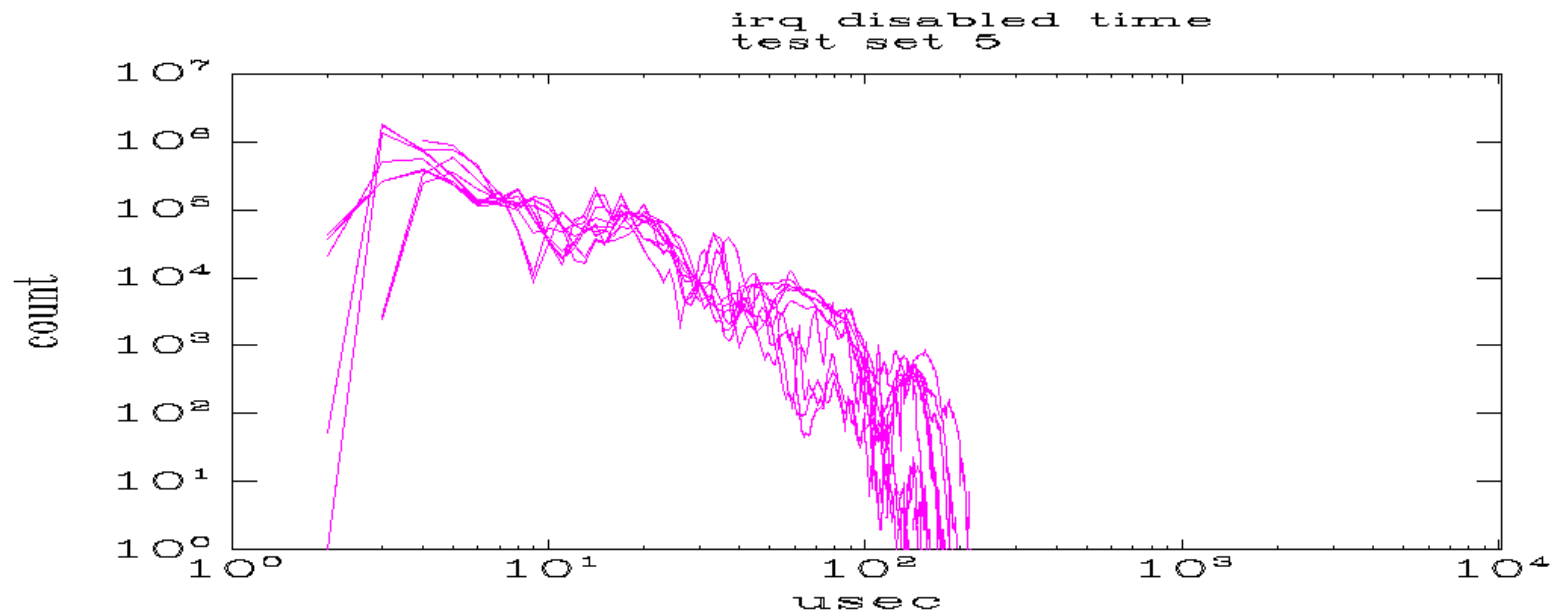


Fix timestamp compare

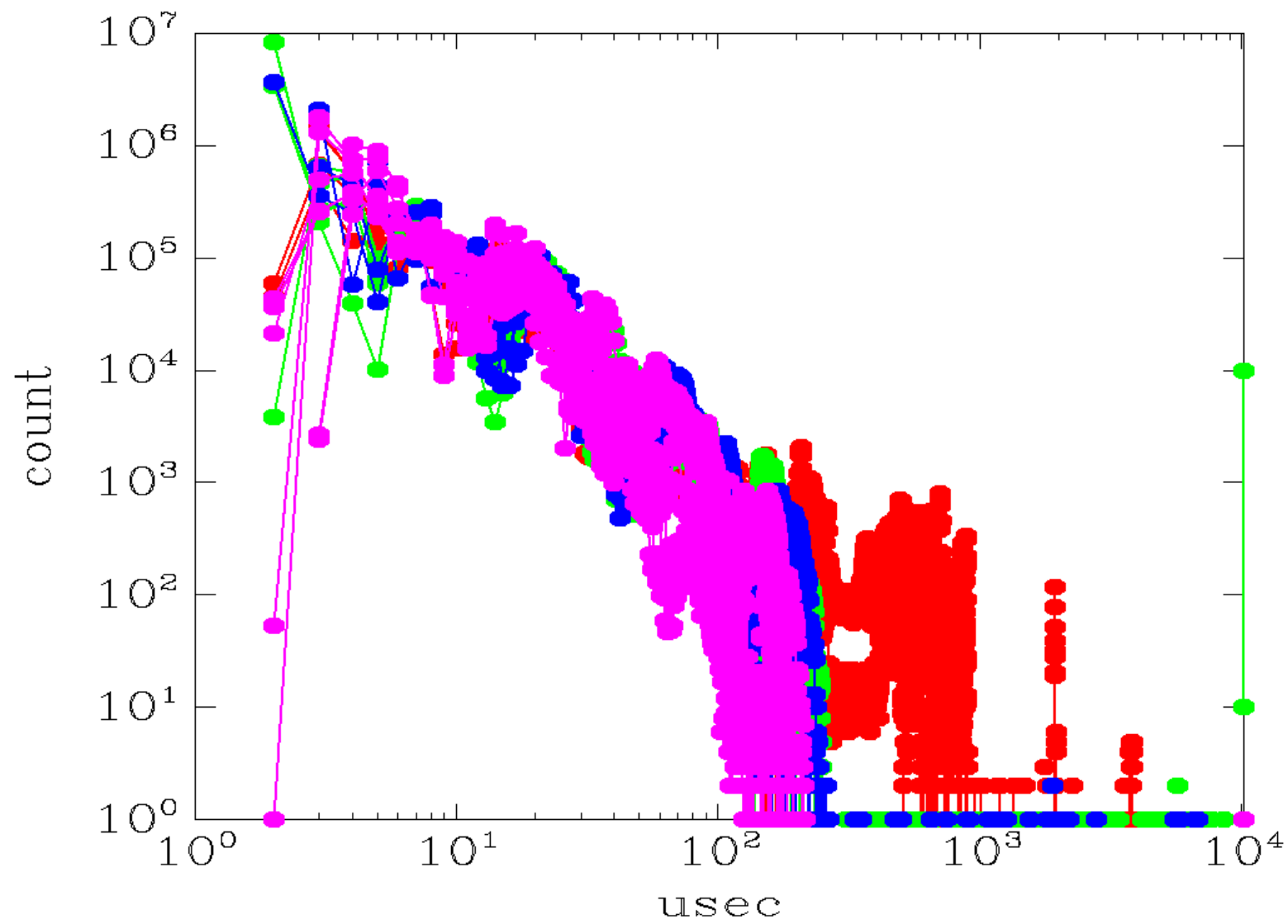




Fix raw / non-raw timestamp transition



irq disabled time
changes from instrumentation fixes



The truly outrageous irq disabled times have been resolved.

What to do next?

The truly outrageous irq disabled times have been resolved.

What to do next?

- Fix the thing that disables preemption for the longest time
- Fix the thing that disables interrupts for the longest time

Preempt Off Latency is one of the key components of RT latency

I will ignore preempt off latency in this talk, but essentially the same methods used to tune irq off latency are used to tune preempt off latency

/proc/latency_trace revisited

One tool that is useful for examining paths with irqs disabled or preemption disabled is the latency trace. (The previous example included start and end points, but contained no details in between.)

The next several slides show

- the raw output of a latency trace
- the raw output of a latency trace, slightly edited
- hand annotations of the latency trace

latency: 164 us, #22/22, CPU#0 | (M:rt VP:0, KP:0, SP:1 HP:1)

| task: softirq-timer/0-7 (uid:0 nice:-5 policy:1 rt_prio:50)

=> started at: handle_int+0x10c/0x120 <80021d2c>

=> ended at: schedule+0xac/0x19c <8025ac70>

```

      _-----=> CPU#
    / _-----=> irq5-off
  | / _-----=> need-resched
  || / _-----=> hardirq/softirq
  ||| / _-----=> preempt-depth
  |||| /
  |||||
cmd  \  pid  |||||  time  |  caller
  \  /      |||||  \  |  /
cyclicte-247  0D...    1us+: handle_int+0x10c/0x120 (<0>)
cyclicte-247  0D.h.   14us+: hrtimer_interrupt+0x9c/0x350 ( 1115 1af13
cyclicte-247  0D.h1   20us+: hrtimer_interrupt+0x164/0x350 ( 1115 1af0
cyclicte-247  0D.h1   49us+: activate_task+0x58/0xa4 <<...>-7> (150 1)
cyclicte-247  0D.h1   53us+: __trace_start_sched_wakeup+0xac/0x19c <<.
cyclicte-247  0DNh1   57us : __trace_start_sched_wakeup+0xac/0x19c <<.
cyclicte-247  0DNh1   58us+: try_to_wake_up+0x1d8/0x1e8 <<...>-7> (150
```

(additional entries not shown)

Magic Decoder Ring

cmdline	pid	c p uDNHx	sym_name+offset/size of function (data)
-----	-----	-----	-----
cyclictc	-247	0D.h1 20us+	hrtimer_interrupt+0x164/0x350 (1115 1af1

d irqs off
D irqs hard off

n need resched delayed
N need resched

H hardirq && softirq
h hardirq
s softirq

x preempt count as hex

xxxx!: time_delta to next entry > 100 usec
xxxx+: time_delta to next entry > 1 usec
xxxx : time_delta to next entry <= 1 usec

The entire trace

```
cycligte-247    OD...    1us+: handle_int+0x10c/0x120 (<0>)
cycligte-247    OD.h.    14us+: hrtimer_interrupt+0x9c/0x350 ( 1115 1af13
cycligte-247    OD.h1    20us+: hrtimer_interrupt+0x164/0x350 ( 1115 1af0
cycligte-247    OD.h1    49us+: activate_task+0x58/0xa4 <<...>-7> (150 1)
cycligte-247    OD.h1    53us+: __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    57us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    58us+: try_to_wake_up+0x1d8/0x1e8 <<...>-7> (150
cycligte-247    ODNh1    87us : activate_task+0x58/0xa4 <<...>-14> (150 2
cycligte-247    ODNh1    88us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    90us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    91us+: try_to_wake_up+0x1d8/0x1e8 <<...>-14> (15
cycligte-247    ODNh1    97us+: enqueue_hrtimer+0x4c/0x1b4 ( 1115 1b2e020
cycligte-247    ODNh.    106us+: clockevents_program_event+0x9c/0x290 ( 11
cycligte-247    ODNh2    128us : activate_task+0x58/0xa4 <<...>-27> (150 3
cycligte-247    ODNh2    130us+: __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh2    132us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh2    133us+: try_to_wake_up+0x1d8/0x1e8 <<...>-27> (15
<...>-7        OD..1    155us+: __schedule+0x3b8/0x638 <cycligte-247> (0 :
<...>-7        OD...    160us+: schedule+0xac/0x19c (<0>)
<...>-7        OD...    162us : trace_hardirqs_on+0xd4/0xf4 (schedule+0xa
```

What are the trace points?

The trace is NOT a list of all functions executed or code paths traversed.

What are the trace points?

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The trace is a collection of interesting locations in various subsystems, which provide a sense of the general flow of control and some related data at some of those locations.

What are the trace points?

The trace is NOT a list of all functions executed or code paths traversed.

The trace is a collection of interesting locations in various subsystems, which provide a sense of the general flow of control and some related data at some of those locations.

It is up to the analyst to interpolate between those locations (and add additional temporary trace points if needed).

Data Fields

(The left part of data has been truncated to show the right half)

```
0D... handle_int+0x10c/0x120 (<0>)
0D.h. hrtimer_interrupt+0x9c/0x350 ( 1115 1af136f1 0)
0D.h1 hrtimer_interrupt+0x164/0x350 ( 1115 1af0f900 80340400)
0D.h1 activate_task+0x58/0xa4 <<...>-7> (150 1)
0D.h1 __trace_start_sched_wakeup+0xac/0x19c <<...>-7> (49 -1)
0DNh1 __trace_start_sched_wakeup+0xac/0x19c <<...>-7> (49 -1)
0DNh1 try_to_wake_up+0x1d8/0x1e8 <<...>-7> (150 0)
0DNh1 activate_task+0x58/0xa4 <<...>-14> (150 2)
0DNh1 __trace_start_sched_wakeup+0xac/0x19c <<...>-14> (49 -1)
0DNh1 __trace_start_sched_wakeup+0xac/0x19c <<...>-14> (49 -1)
0DNh1 try_to_wake_up+0x1d8/0x1e8 <<...>-14> (150 0)
0DNh1 enqueue_hrtimer+0x4c/0x1b4 ( 1115 1b2e0200 80340400)
0DNh. clockevents_program_event+0x9c/0x290 ( 1115 1af42931 99376)
0DNh2 activate_task+0x58/0xa4 <<...>-27> (150 3)
0DNh2 __trace_start_sched_wakeup+0xac/0x19c <<...>-27> (49 -1)
0DNh2 __trace_start_sched_wakeup+0xac/0x19c <<...>-27> (49 -1)
0DNh2 try_to_wake_up+0x1d8/0x1e8 <<...>-27> (150 0)
0D..1 __schedule+0x3b8/0x638 <cyclicte-247> (0 150)
0D... schedule+0xac/0x19c (<0>)
0D... trace_hardirqs_on+0xd4/0xf4 (schedule+0xac/0x19c)
```

Data Fields Hand Annotated

(no Magic Decoder Ring)

```
latency: 164 us, #22/22, CPU#0 | (M:rt VP:0, KP:0, SP:1 HP:1)
=> started at: handle_int+0x10c/0x120 <80021d2c>
=> ended at:   schedule+0xac/0x19c <8025ac70>
```

```
handle_int+0x10c/0x120 (<0>)
hrtimer_interrupt+0x9c/0x350 ( 1115 1af136f1 0)
hrtimer_interrupt+0x164/0x350 ( 1115 1af0f900 80340400)
activate_task+0x58/0xa4 <<...>-7> (150 1)
__trace_start_sched_wakeup+0xac/0x19c <<...>-7> (49 -1)
__trace_start_sched_wakeup+0xac/0x19c <<...>-7> (49 -1)
try_to_wake_up+0x1d8/0x1e8 <<...>-7> (150 0)
activate_task+0x58/0xa4 <<...>-14> (150 2)
__trace_start_sched_wakeup+0xac/0x19c <<...>-14> (49 -1)
__trace_start_sched_wakeup+0xac/0x19c <<...>-14> (49 -1)
try_to_wake_up+0x1d8/0x1e8 <<...>-14> (150 0)
enqueue_hrtimer+0x4c/0x1b4 ( 1115 1b2e0200 80340400)
clockevents_program_event+0x9c/0x290 ( 1115 1af42931 99376)
activate_task+0x58/0xa4 <<...>-27> (150 3)
__trace_start_sched_wakeup+0xac/0x19c <<...>-27> (49 -1)
__trace_start_sched_wakeup+0xac/0x19c <<...>-27> (49 -1)
try_to_wake_up+0x1d8/0x1e8 <<...>-27> (150 0)
__schedule+0x3b8/0x638 <cyclicte-247> (0 150)
schedule+0xac/0x19c (<0>)
trace_hardirqs_on+0xd4/0xf4 (schedule+0xac/0x19c)
```

```
time 1115 1af136f1
time 1115 1af0f900 timer 80340400
pid 7, PRIO 150, nr_running 1
pid 7, prio 49
pid 7, prio 49
pid 7, PRIO 150, PRIO(rq->curr) 0
pid 14, PRIO 150, nr_running 2
pid 14, prio 49
pid 14, prio 49
pid 14, PRIO 150, PRIO(rq->curr) 0
time 1115 1b2e0200 timer 80340400
time 1115 1af42931, delta 99376
pid 27, PRIO 150, nr_running 3
pid 27, prio 49
pid 27, prio 49
pid 27, PRIO 150, PRIO(rq->curr) 0
pid 247, PRIO was 0, PRIO is 150
```

Finding source location, data fields (the easy way)

```
OD.h.  hrtimer_interrupt+0x9c/0x350 ( 1115 1af136f1 0)
OD.h1  hrtimer_interrupt+0x164/0x350 ( 1115 1af0f900 80340400)
```

Finding source location, data fields (the easy way)

```
OD.h.  hrtimer_interrupt+0x9c/0x350 ( 1115 1af136f1 0)
OD.h1  hrtimer_interrupt+0x164/0x350 ( 1115 1af0f900 80340400)
```

Look at the source, maybe it's obvious
(or maybe it's not....).

```
void hrtimer_interrupt(struct clock_event_device *dev)
{
    .....
    retry:
        now = ktime_get();
        hrtimer_trace(now, 0);
        .....

        hrtimer_trace(timer->expires, (unsigned long) timer);
}
```

Finding source location, data fields

The hard way, using gdb.

Finding source location, data fields

(gdb) i line *hrtimer_interrupt+0x9c

Line 1105 of "kernel/hrtimer.c"

starts at address 0x8006ecf4 <hrtimer_interrupt+156>

ends at 0x8006ed18 <hrtimer_interrupt+192>.

1102   now = ktime_get();

1103   hrtimer_trace(now, 0);

1104

1105   expires_next.tv64 = KTIME_MAX;

Finding source location, data fields

(gdb) i line *hrtimer_interrupt+0x9c

Line 1105 of "kernel/hrtimer.c"

starts at address 0x8006ecf4 <hrtimer_interrupt+156>

ends at 0x8006ed18 <hrtimer_interrupt+192>.

```
1102         now = ktime_get();
1103         hrtimer_trace(now, 0);
1104
1105         expires_next.tv64 = KTIME_MAX;
```

Note that the trace address is usually the location the trace function returns to.

```
# define hrtimer_trace(a,b)    trace_special((a).tv.sec,(a).tv.nsec,b)
```

Finding source location, data fields

... a more ugly example

(gdb) i line *schedule+0xac

Line 43 of "irqflags.h" starts at address 0x8025ac70 <schedule+172>
and ends at 0x8025ac90 <schedule+204>.

```
41 static inline void raw_local_irq_enable(void)
42 {
43     __asm__ __volatile__(
        "        mfc0    $1,$12 \n"
```

(gdb) i line *schedule+0xac

Line 43 of "irqflags.h" starts at address 0x8025ac70 <schedule+172> and ends at 0x8025ac90 <schedule+204>.

```
41 static inline void raw_local_irq_enable(void)
42 {
43     __asm__ __volatile__(
        "        mfc0    $1,$12 \n"
```

(gdb) x/i *schedule+0xac

0x8025ac70 <schedule+172>: mfc0 at,\$12

(gdb) x/8i *schedule+0xac - 0x10

```
0x8025ac60 <schedule+156>:    bnez     v0,0x8025ac38 <schedule+116>
0x8025ac64 <schedule+160>:    nop
0x8025ac68 <schedule+164>:    jal     0x80084c54 <trace_hardirqs_on>
0x8025ac6c <schedule+168>:    nop
0x8025ac70 <schedule+172>:    mfc0    at,$12
0x8025ac74 <schedule+176>:    nop
0x8025ac78 <schedule+180>:    ori     at,at,0x1f
0x8025ac7c <schedule+184>:    xori    at,at,0x1e
```

(gdb) i line ***schedule+0xac**

Line 43 of "irqflags.h" starts at address 0x8025ac70 <schedule+172>
and ends at 0x8025ac90 <schedule+204>.

(gdb) x/i ***schedule+0xac - 0x10**

0x8025ac60 <schedule+156>:   bnez   v0,0x8025ac38 <schedule+116>

(gdb) i line ***0x8025ac60**

Line 3854 of "kernel/sched.c"

starts at address 0x8025ac60 <schedule+156>

ends at 0x8025ac68 <schedule+164>.

```
3852        do {
3853            __schedule();
3854        } while (unlikely(test_thread_flag(TIF_NEED_RESCHED) ||
3855                          test_thread_flag(TIF_NEED_RESCHED_DELAYED)));
3856
3857        local_irq_enable();
```

The entire trace

```
cycligte-247    OD...    1us+: handle_int+0x10c/0x120 (<0>)
cycligte-247    OD.h.    14us+: hrtimer_interrupt+0x9c/0x350 ( 1115 1af13
cycligte-247    OD.h1    20us+: hrtimer_interrupt+0x164/0x350 ( 1115 1af0
cycligte-247    OD.h1    49us+: activate_task+0x58/0xa4 <<...>-7> (150 1)
cycligte-247    OD.h1    53us+: __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    57us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    58us+: try_to_wake_up+0x1d8/0x1e8 <<...>-7> (150
cycligte-247    ODNh1    87us : activate_task+0x58/0xa4 <<...>-14> (150 2
cycligte-247    ODNh1    88us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    90us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh1    91us+: try_to_wake_up+0x1d8/0x1e8 <<...>-14> (15
cycligte-247    ODNh1    97us+: enqueue_hrtimer+0x4c/0x1b4 ( 1115 1b2e020
cycligte-247    ODNh.    106us+: clockevents_program_event+0x9c/0x290 ( 11
cycligte-247    ODNh2    128us : activate_task+0x58/0xa4 <<...>-27> (150 3
cycligte-247    ODNh2    130us+: __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh2    132us : __trace_start_sched_wakeup+0xac/0x19c <<.
cycligte-247    ODNh2    133us+: try_to_wake_up+0x1d8/0x1e8 <<...>-27> (15
<...>-7        OD..1    155us+: __schedule+0x3b8/0x638 <cycligte-247> (0 :
<...>-7        OD...    160us+: schedule+0xac/0x19c (<0>)
<...>-7        OD...    162us : trace_hardirqs_on+0xd4/0xf4 (schedule+0xa
```

Leads to Interesting Kernel Path

1) Timer interrupt

- highres timers code
- wake appropriate threads
- schedule

Leads to Interesting Kernel Path

1) Timer interrupt

- highres timers code
- wake appropriate threads
- schedule

$O(n)$ algorithm -- more timers expiring at same time will result in a longer maximum IRQ off

Obvious in the latency_trace we were examining.

Possible Fix

Not investigated yet.

Possible Workaround

Avoid large number of timers expiring at the same time.

Interesting Kernel Path

2) interrupt top half handling followed by `preempt_schedule_irq()` is a long path with irqs disabled

Found by looking at the intermediate time stamps in a latency trace.

Possible Workaround

Remove or rate limit non-RT interrupts.

Possible Workaround

Remove or rate limit non-RT interrupts.

In my case, the large interrupt volume is due to network traffic since my root file system is NFS mounted from another host. It is not realistic for any use case that I am expecting.

Possible Fix

resume_kernel:

THIS IS NOT RECOMMENDED, do not do this unless you really
understand the negative effects of enabling irqs here

raw_local_irq_enable t0
raw_local_irq_disable

lw t0, kernel_preemption
beqz t0, restore_all
lw t0, TI_PRE_COUNT(\$28)
bnez t0, restore_all

need_resched:

<< code deleted for brevity >>
jal preempt_schedule_irq

WARNING

Enabling irqs on the return from interrupts path allows nested interrupts, which may result in a stack overflow.

If you do not understand the negative effects of allowing nested interrupts, or can not ensure they will not crash your specific system, do not apply this change.

The next two slides are tx49 irq disabled time for two cases:

red:

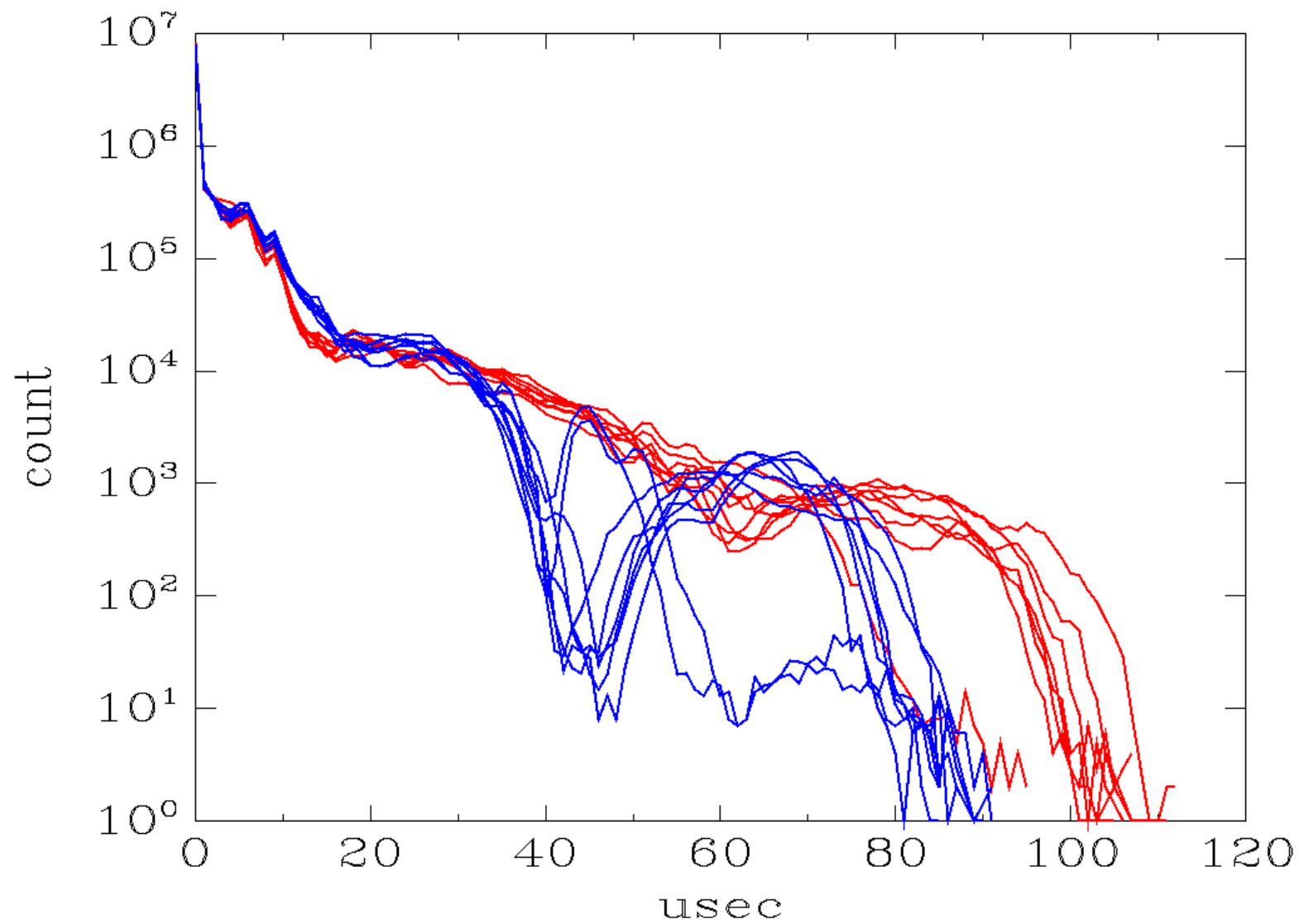
baseline

blue:

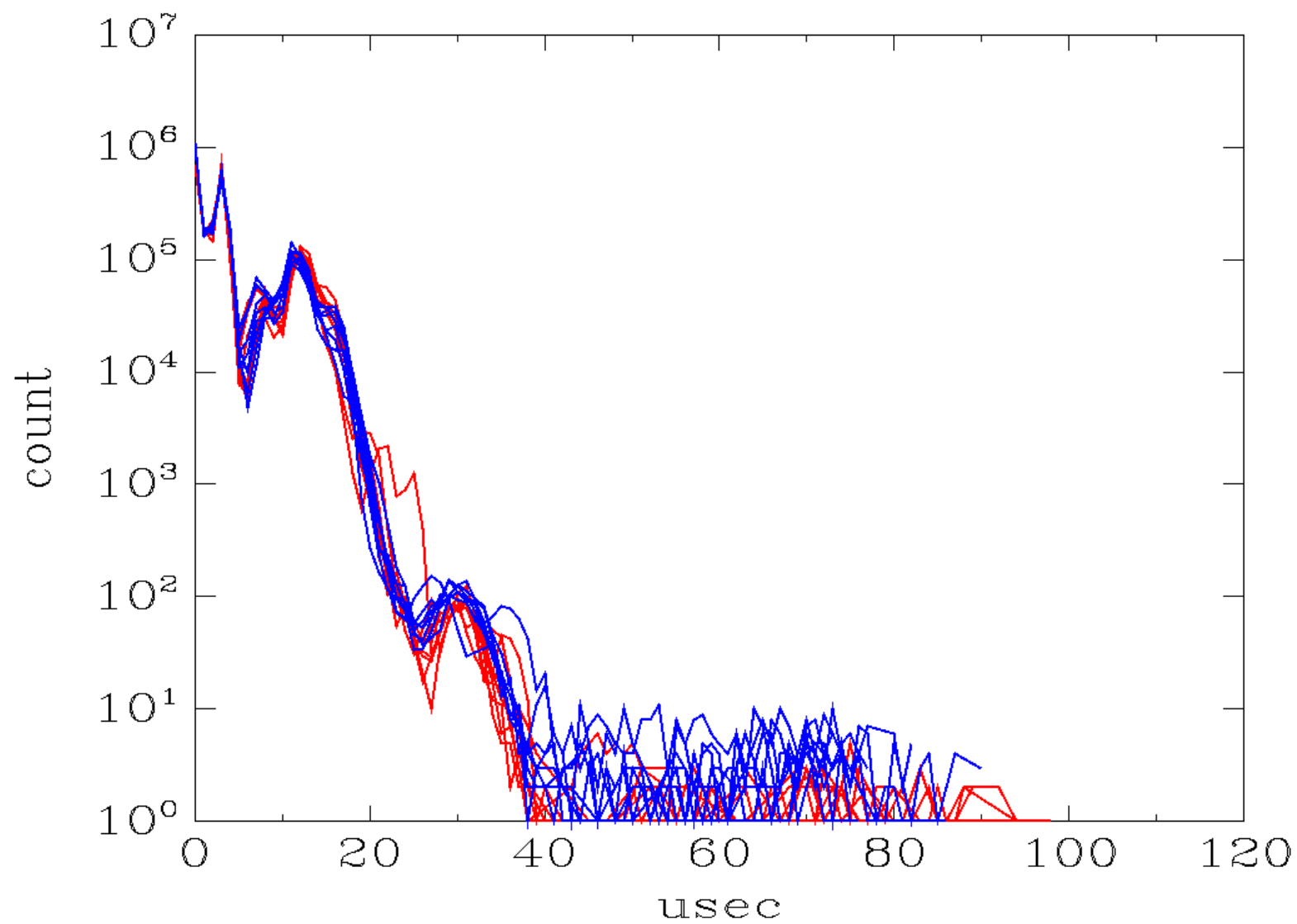
Fix for Interesting Kernel Path 2
in resume_kernel

Good improvement for irq off time.

irq disabled time   heavy load
change from enabling irq in resume_kernel



irq disabled time   no load
change from enabling irqs in resume_kernel



before resume_kernel fix

with resume_kernel fix, allow only rt irqs in break disable

Max irq off

112 usec heavy load

91 usec heavy load

100 usec no load

86 usec no load

cyclictest max wakeup latency

180 usec heavy load

269 usec heavy load

139 usec no load

245 usec no load

The Moral

Do not lose sight of the most important metric -- meeting the real time application deadline -- while trying to tune the components that cause latency.

The Real Solution

Modify the algorithms and/or data structures to shorten the code path. Probably a major project.

Is this fix reasonable?

No, not really, except in extremely constrained cases. You will never see me submit these simplistic patches upstream. They are not appropriate for a general purpose real time operating system.

Potentially useful if all of the real-time processing is occurring in kernel drivers.

The code shown is much more simplistic than the complete solution that might be reasonable.

Yet another tool

LatencyTOP

Development release 0.1 on Jan 18, 2008

www.latencytop.org

The next tuning tool

Tap into the experts' knowledge -- the web is your friend!

Search engines, wikis, web sites, email lists...

The “Resources” slides at the end of this presentation references some good sources of information.

Normal mail list etiquette applies

- 1) Search the history
  - The issue and a solution for it may be known.
- 2) Try to solve the problem yourself
- 3) Then, consider asking on the list
  - Do your homework, present the data in a concise but complete form.
  - Be prepared to provide additional data and clarifications, collect additional data, and test suggested solutions upon request.
  - Don't expect other people to do your work.
  - Etc...

An example from linux-rt-user

Latencies up to 600us for

- 2.6.24-rc8-rt1: mpc5200 powerpc
- NFS mounted root filesystem
- CONFIG_PREEMPT_RCU_BOOST
or CONFIG_RCU_TRACE not set
or CONFIG_RCU_TRACE=m

Start of the thread discussing this is on LKML
and linux-rt-user:

Subject: Re: 2.6.24-rc8-rt1
From: Wolfgang Grandegger
Date: Thu Jan 17 2008

Thread subject morphs into:

Re: 2.6.24-rc8-rt1: Strange latencies on
mpc5200 powerpc

Some examples of recent history

Softirq processing

Are there any known issues, will they be fixed?

USB subsystem (Isochronous might be OK).

Is it usable for a real-time project?

From the mpc5200 powerpc latency thread:

“It's also my suspicion that the high latencies are related to the RCU usage in the network layer, where it's heavily used.”

Current performance results

Very “preliminary” since tuning has not been completed, but there are some things that can be said.

Kernel version:

MIPS 2.6.24 + patch-2.6.24-rt1 + tx4937 fixes

Source of the data on the following slides

(1) “Realtime capabilities of low-end PowerPC
and ARM boards for embedded systems”

Alexander Bauer

9th Real Time Linux Workshop

(2, 3, 4, 5) Frank Rowand, March 2008

***** warning: comparing unlike metrics *****

ppc MCP 5200	266 Mhz	120 usec	(1)
tx49	250 Mhz	90 usec	(2)
tx49	250 Mhz	139 usec	(3)
tx49	250 Mhz	91 usec	(4)
tx49	250 Mhz	180 usec	(5)
arm PXA270	260 Mhz	600 usec	(1)

(1) moderate load (3 cyclicttest threads, ping flood)

metric: cyclicttest max wakeup latency

(2,3) light load (5 cyclicttest threads)

(4,5) heavy load (5 cyclicttest threads, 9 ls -IR)

0% cpu idle

(2,4) metric: max IRQ disabled time

(3,5) metric: cyclicttest max wakeup latency

Alexander Bauer Tests

ppc MCP 5200
arm PXA270

cyclicttest -q -n -t 3 -p 99

Background load:

ping flood

Frank Rowand Tests

Toshiba TX4937 Reference Board

`cyclictest -p 80 -t 5 -n -l 100000`

Background load either:

1) none

2) 9 instances of:

`ls -lR / >/dev/null`

/ is nfs mounted from host

Frank Rowand Tests - priorities

PID	COMMAND	RTPRIO	CLS
3	IRQ-7	50	FF
4	IRQ-11	50	FF
5	posix_cpu_timer	99	FF
6	softirq-high/0	50	FF
7	softirq-timer/0	50	FF
8	softirq-net-tx/	50	FF
9	softirq-net-rx/	50	FF
10	softirq-block/0	50	FF
11	softirq-tasklet	50	FF
12	softirq-sched/0	50	FF
13	softirq-hrtimer	50	FF
14	softirq-rcu/0	50	FF
16	events/0	1	FF
19	krcupreemptd	1	FF
24	IRQ-13	50	FF
26	IRQ-16	50	FF
137	cyclictest	-	TS
138	cyclictest	-	TS
140	cyclictest	80	FF
142	cyclictest	79	FF
143	cyclictest	78	FF
144	cyclictest	77	FF
145	cyclictest	76	FF

Measurement Overhead

Enabling measurement instrumentation adds significant overhead. Remember to disable it before trying to measure actual real time behaviour.

Measurement Overhead

Console boot messages will remind you:

```
Looking up port of RPC 100003/2 on 192.168.1.10
```

```
Looking up port of RPC 100005/1 on 192.168.1.10
```

```
VFS: Mounted root (nfs filesystem).
```

```
*****
```

```
*                                                                 *
```

```
*  REMINDER, the following debugging options are turned on in your .config: *
```

```
*                                                                 *
```

```
*      CONFIG_CRITICAL_PREEMPT_TIMING                               *
```

```
*      CONFIG_CRITICAL_IRQSOFF_TIMING                             *
```

```
*      CONFIG_FUNCTION_TRACE                                       *
```

```
*                                                                 *
```

```
*  they may increase runtime overhead and latencies.              *
```

```
*                                                                 *
```

```
*****
```

```
Freeing init memory: 116K
```

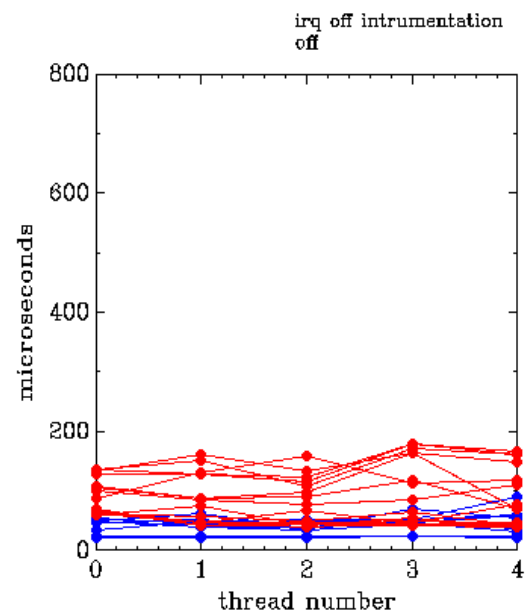
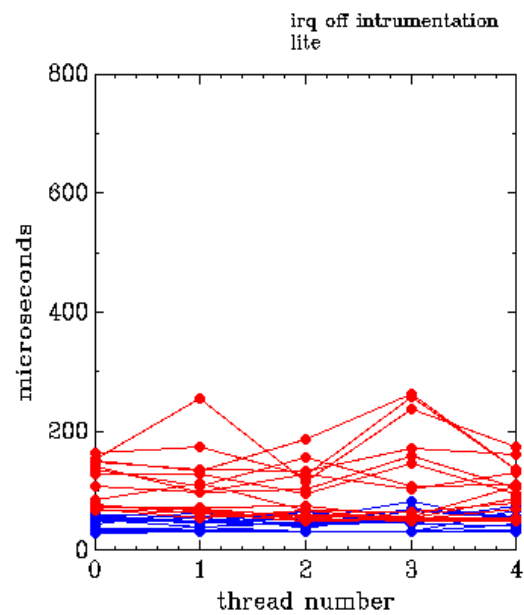
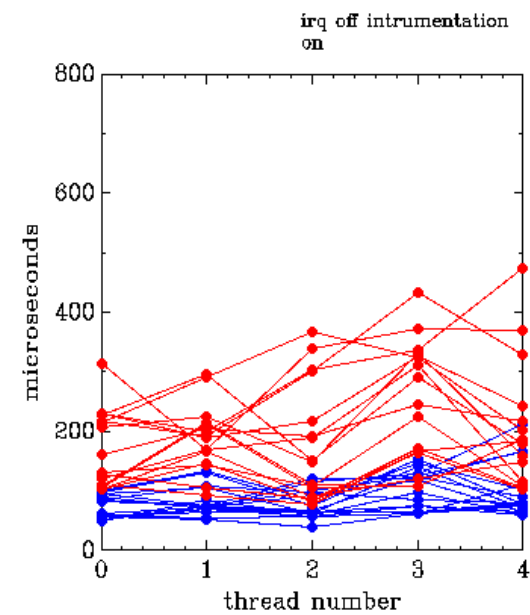
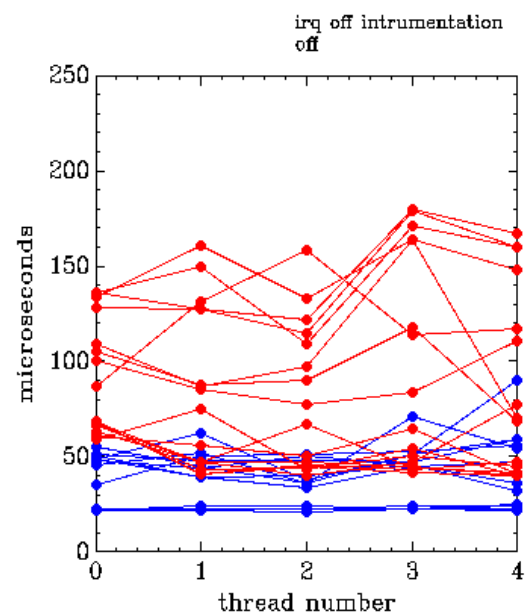
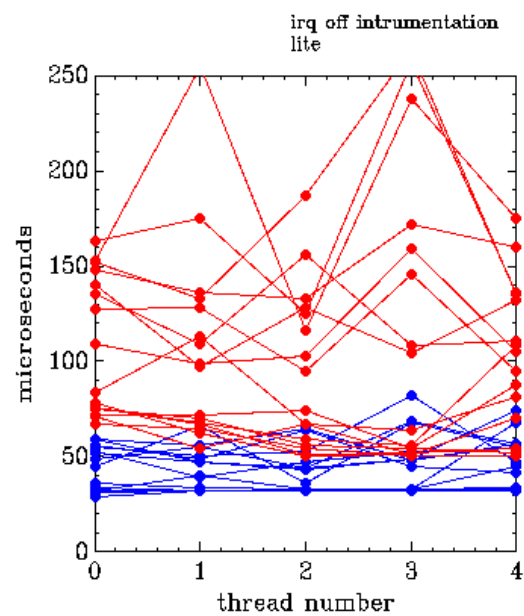
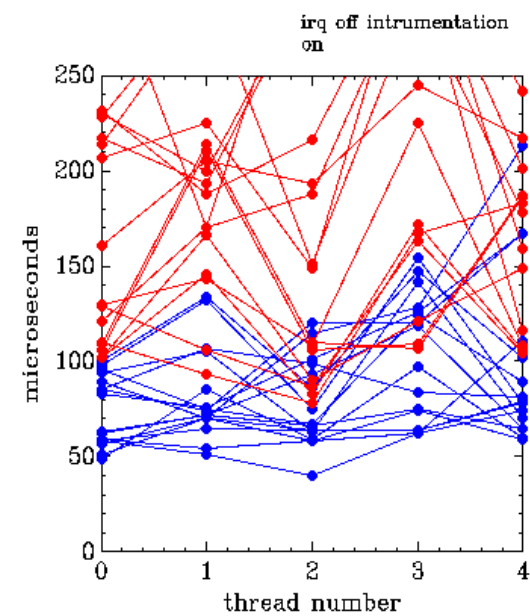
The next slide is tx49 cyclicttest latency data (blue is average, red is max), for three cases:

- no latency tracing enabled
- trace-lite enable
- preempt-rt patch latency tracing enabled

The data shows the large performance impact of enabling latency tracing.

Each individual graph contains multiple lines, where each line is the result of a single test run.

1/2 of the tests have no “ls” background load.



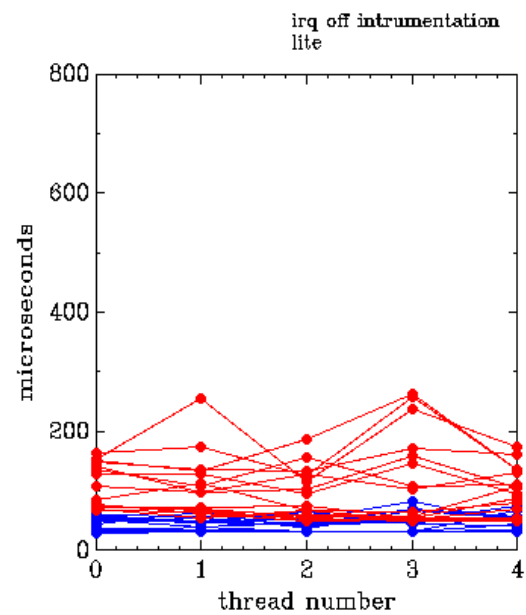
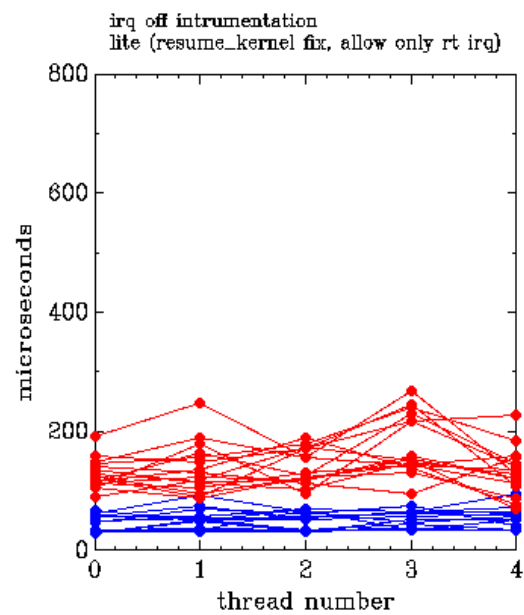
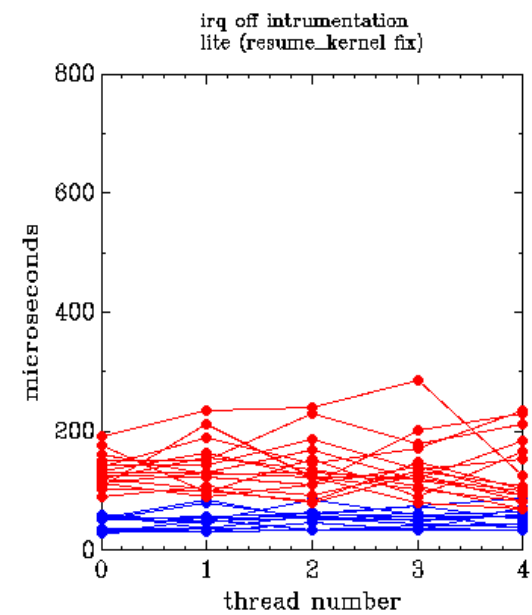
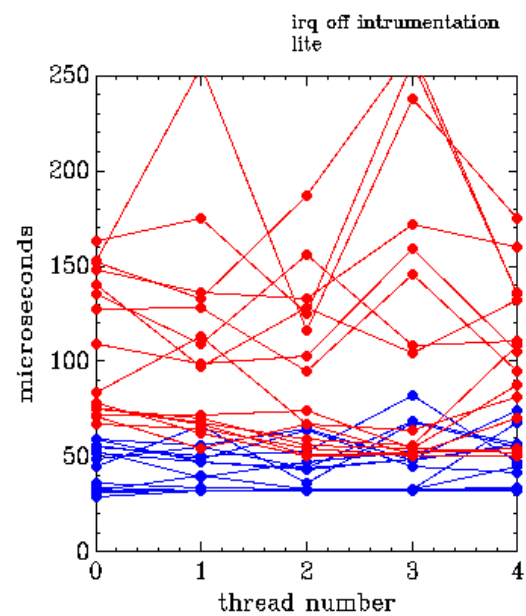
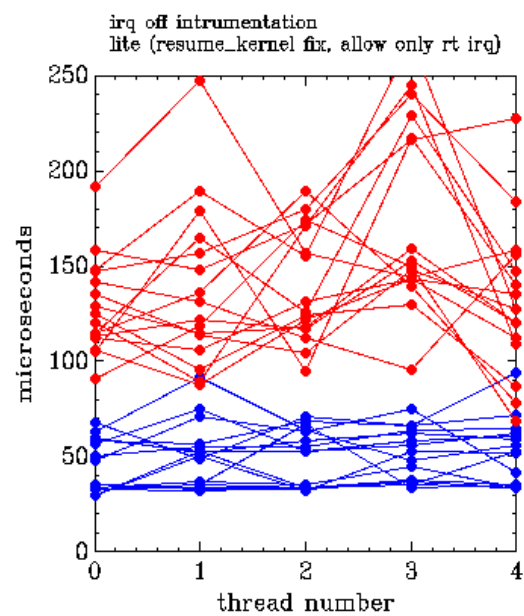
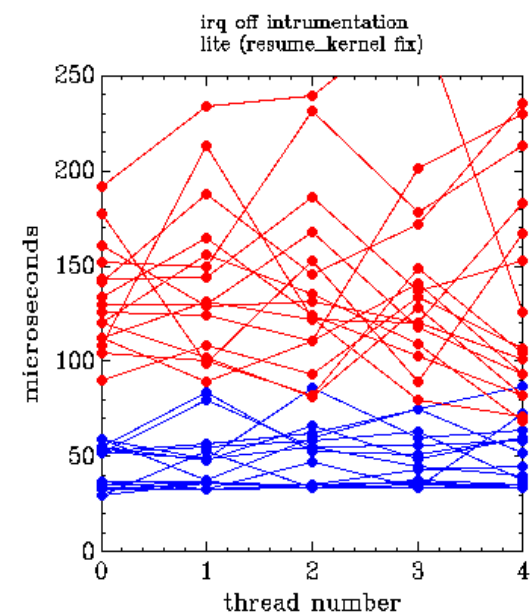
The next slide is tx49 cyclicttest latency for three cases:

- baseline
- resume_kernel fix, allow only timer irqs in break
- resume_kernel fix

Good improvement was seen earlier for irqs off time.

But worst case cyclicttest results suffer.

This is an example of how tuning for a single metric can harm the overall RT application.



The next slides are tx49 interrupts disabled time for two cases:

red:

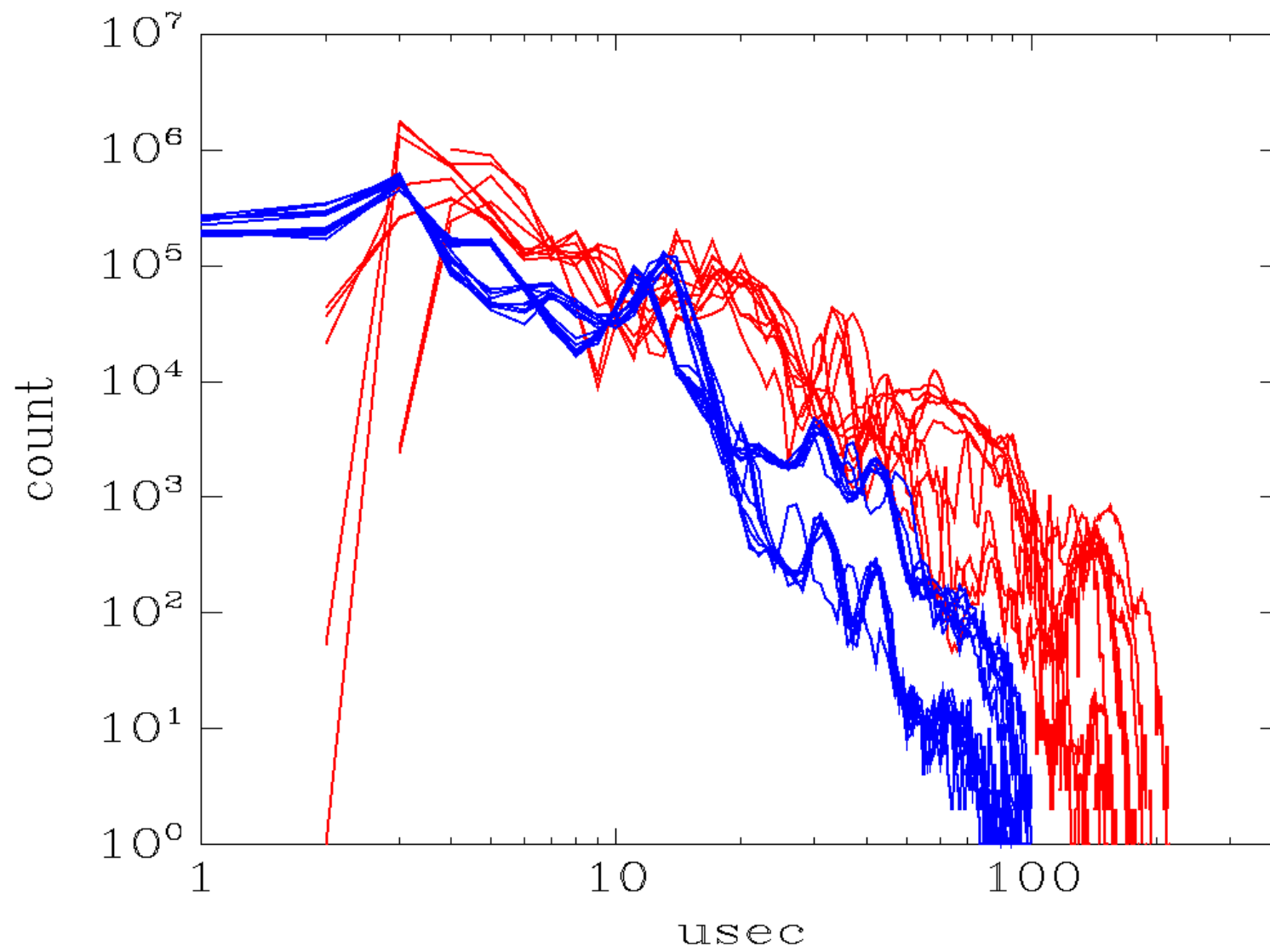
preempt-rt patch latency tracing enabled

blue:

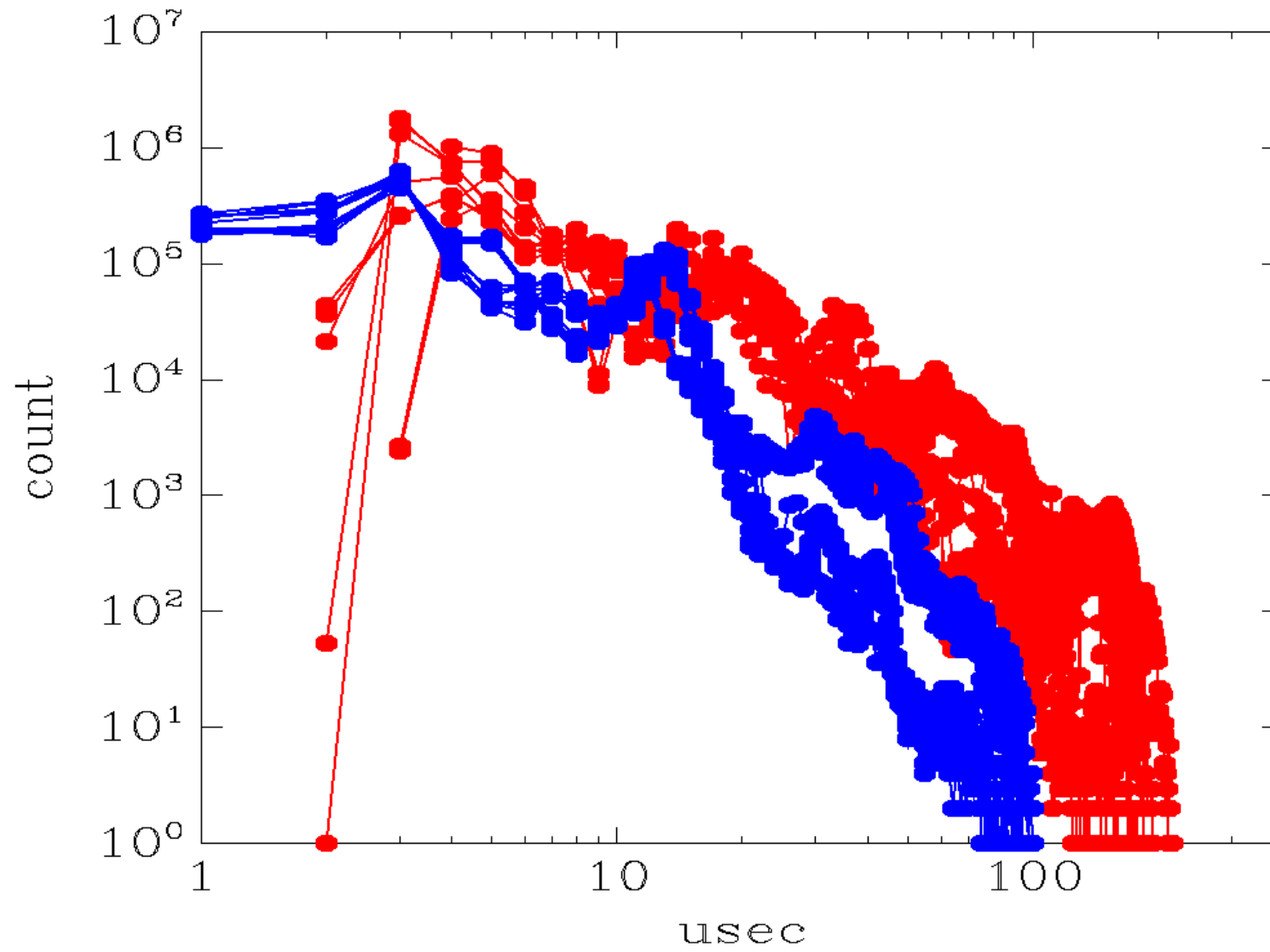
trace-lite enabled

The data again shows the large performance impact of enabling latency tracing.

irq disabled time
shows overhead of instrumentation



irq disabled time
shows overhead of instrumentation



Imbench “results”

The overhead for non-basic operations on RT-PREEMPT varies from small to moderate to large.

A few cases where the rt-preempt overhead is smaller than for SMP kernel on UP hardware.

There are inconsistencies between test runs (3 runs per configuration) which implies the data is not reliable.

No attempt was made to validate this data.

Is this Imbench data valid?

It should be viewed with suspicion.

No attempt was made to validate the results (this was just a quick attempt to collect a few data points to see if they would provide some insights into the overhead of preempt-rt).

Not consistent with other reports of recent versions of the kernel on x86, such as Siro Arthur, et al, 9th Real Time Linux Workshop.

Is this Imbench data valid?

Siro Arthur, et al

“Apparently PREEMPT RT has no significant degrading impact on the general performance of the system in its current version”
(2.6.21.5, 2.6.22.1, 2.6.23-rc1)

[for] “...tests against previous version[s] e.g. 2.6.14-rt20 ... the performance of these kernels were significantly below that of the unpatched versions”
(apologies to Siro for severely mangling this quote)

Limitations of the test methodology

- Non-representative workload.
- No attempt to exercise all areas of system functionality.
- Extremely short test duration.

Bottom line:

This is just the start of this specific tuning project.

The new tracer, 'ftrace', is in rt patchset 2.6.24-rt2 mainline 2.6.27-rc1

old file name

/proc/latency_trace

/proc/latency_hist/interrupt_off_latency/CPU0
/proc/latency_hist/preempt_off_latency/CPU0
/proc/latency_hist/wakeup_latency/CPU0

new file name

/debugfs/tracing/latency_trace
/debugfs/tracing/trace

/debugfs/tracing/
/debugfs/tracing/
/debugfs/tracing/

see Documentation/ftrace.txt for more information

LKML:

From: Ingo Molnar

Date: Sun Feb 10 2008

Subject: [10/19] ftrace: add basic support for gcc profiler instrumentation

Subject: [11/19] ftrace: latency tracer infrastructure for documentation

Subject: [12/19] ftrace: function tracer

Subject: [13/19] ftrace: add tracing of context switches

Subject: [14/19] ftrace: tracer for scheduler wakeup latency

Subject: [15/19] ftrace: trace irq disabled critical timings

Subject: [16/19] ftrace: trace preempt off critical timings

Resources

Rtiwiki

http://rt.wiki.kernel.org/index.php/Main_Page

rt-user-list

<http://dir.gmane.org/gmane.linux.rt.user>

eLinux.org

http://elinux.org/Real_Time

cyclictest

<http://git.kernel.org/?p=linux/kernel/git/tglx/rt-tests.git;a=summary>

Resources

ftrace

<http://people.redhat.com/srostedt/ftrace-tutorial.odp>

kernel source: Documentation/ftrace.txt

hackbench

<http://devresources.linux-foundation.org/craiger/hackbench/>

LatencyTOP

<http://www.latencytop.org>

“Stress actions - things that will kill realtime performance”
and information about test programs and testing

http://elinux.org/Realtime_Testing_Best_Practices

Resources

A realtime preemption overview

<http://lwn.net/Articles/146861>

What's in the realtime tree

<http://lwn.net/Articles/252716>

Ninth Real-Time Linux Workshop 2007

<http://lwn.net/Articles/260118>

<http://linuxdevices.com/articles/AT4991083271.html>